

## Homework 6 — ML module language — assigned Monday 17 April — due Tuesday 25 April

All code in this homework assignment must use the SML module language, and it must be organized using the SML/NJ Compilation Manager. Place all functor applications together in a file `link.sml`. Put all tests into a structure `Tests` in the file `tests.sml`.

### 6.1 Lambda-calculus (30pts)

#### A library of $\lambda$ -terms

$$\begin{aligned}
 \mathbf{I} &\triangleq \lambda x.x & \mathbf{K} &\triangleq \lambda xy.x & \mathbf{S} &\triangleq \lambda fgx.(fx)(gx) & \mathbf{B} &\triangleq \lambda fgx.f(gx) & \mathbf{C} &\triangleq \lambda fgx.fxg \\
 \omega &\triangleq \lambda x.xx & \Omega &\triangleq \omega\omega & \mathbf{Y} &\triangleq \lambda f.(\lambda x.f(xx))(\lambda x.f(xx)) \\
 \mathbf{true} &\triangleq \lambda xy.x & \mathbf{false} &\triangleq \lambda xy.y & \mathbf{not} &\triangleq \lambda t.t \mathbf{false} \mathbf{true} & \mathbf{cond} &\triangleq \lambda ee_1e_2.ee_1e_2 \\
 \mathbf{pair} &\triangleq \lambda e_1e_2f.f e_1e_2 & \mathbf{fst} &\triangleq \lambda p.p \mathbf{true} & \mathbf{snd} &\triangleq \lambda p.p \mathbf{false} \\
 \mathbf{0} &\triangleq \lambda fx.x & \mathbf{1} &\triangleq \lambda fx.fx & \mathbf{2} &\triangleq \lambda fx.f(fx) & \mathbf{succ} &\triangleq \lambda nfx.nf(fx) & \mathbf{add} &\triangleq \lambda mnfx.mf(nfx) \\
 \mathbf{iszero} &\triangleq \lambda n.n(\lambda x.\mathbf{false})\mathbf{true} & \mathbf{prefn} &\triangleq \lambda fp.\mathbf{pair} \mathbf{false}(\mathbf{cond}(\mathbf{fst} p)(\mathbf{snd} p)(f(\mathbf{snd} p))) \\
 \mathbf{pred} &\triangleq \lambda nfx.\mathbf{snd}(n(\mathbf{prefn} f)(\mathbf{pair} \mathbf{true} x)) \\
 \mathbf{cons} &\triangleq \lambda hts.sht & \mathbf{hd} &\triangleq \lambda L.L \mathbf{true} & \mathbf{tl} &\triangleq \lambda L.L \mathbf{false} & \mathbf{nil} &\triangleq \lambda x.\mathbf{true} & \mathbf{isempty} &\triangleq \lambda L.L(\lambda ht.\mathbf{false})
 \end{aligned}$$

#### Normal forms of some $\lambda$ -terms

$$\begin{aligned}
 \mathbf{SKK} &\rightarrow \lambda x.x & \mathbf{K(SII)} &\rightarrow \lambda ab.bb & \mathbf{S(S(KS)(KI))(KI)} &\rightarrow \lambda ab.bb \\
 \mathbf{SSSSSSS} &\rightarrow \lambda ab.(ab(ab(ab\lambda c.ac(bc))))
 \end{aligned}$$

#### 6.1.1

Show that the following  $\lambda$ -terms have a normal form:

1.  $(\lambda y.yyy)((\lambda ab.a)\mathbf{I}(\mathbf{SS}))$
2.  $(\lambda yz.zy)((\lambda x.xxx)(\lambda x.xxx))(\lambda w.\mathbf{I})$

#### 6.1.2

For each of the following  $\lambda$ -terms either find its normal form or show that it has no normal form:

1.  $(\lambda x.xx)(\lambda x.x)$
2.  $(\lambda x.xx)(\lambda x.xx)$
3.  $\mathbf{Y}$
4.  $\mathbf{Y}(\lambda y.y)$

**6.1.3**

(Turing) Let  $A \triangleq \lambda xy.y(xxy)$ . Let  $\Theta \triangleq AA$ . Show that  $\Theta$  is a fixed-point operator.

**6.2 Lambda-calculus Interpreter (70pts) [K.1.1; K.3.1; K.3.2]**

Develop an interpreter for the  $\lambda$ -calculus that will automate reductions. This program will follow literally the rules for  $\beta$ -conversion and the rules for substitution. The internal representation of  $\lambda$ -terms is essentially the same as the textual representation, though the data type makes the bracketing structure apparent, and pattern-matching easier.

We must first specify the internal representation for  $\lambda$ -terms. The following type must be used:

```
type var = string
datatype expr = Var of var
              | Abs of var * expr
              | App of expr * expr
```

The following tasks build the interpreter bottom-up.

**6.2.1**

Implement an environment mapping identifiers to  $\lambda$ -terms, with type `string -> expr`. There should be a mechanism to build new environments out of old ones by introducing a new definition for an identifier.

**6.2.2**

Implement a function `freeVariables: expr -> var list`.

**6.2.3**

Implement a function `isFreeVariable: var * expr -> bool`.

**6.2.4**

Implement a function `substitute: expr * var * expr -> expr`, such that `substitute (e, x, t)` substitutes  $t$  for  $x$  in  $e$ . To generate fresh variable symbols, use the following code:

```
local
  val counter = ref 0
in
  fun genSym () =
    let
      val x = !counter
```

```

    val _ = counter := x+1
  in
    "_" ^ Int.toString x
  end
end
end

```

### 6.2.5

Implement a function `isBetaRedex: expr -> bool`.

### 6.2.6

Implement a function `convertBetaRedex: expr -> expr`.

### 6.2.7

Implement a function `convert: expr -> expr option` which finds a leftmost outermost  $\beta$  redex, if any, and performs  $\beta$  conversion.

### 6.2.8

Implement a function `reduce: expr -> expr` that applies  $\beta$  conversion steps in normal order until a normal form is found.

### 6.2.9

Test your program by reducing various  $\lambda$ -terms, such as: **S****K****K**; **K**(**S****I**); **S**(**S**(**K****S**)(**K****I**))(**K****I**); **S****S****S****S****S****S**.

### 6.2.10

Implement the factorial function over Church numerals. (Use the **Y** combinator.) Test your program by having it compute  $n!$  for various  $n$ . Report how fast the evaluator works for different inputs or input sizes. (Take into account that with a unary representation, different numbers have different sizes.)

## How to turn in

Submission instructions have been posted to the course mailing list.

Include the following statement with your submission, signed and dated:

*I pledge my honor that in the preparation of this assignment I have complied with the University of New Mexico Board of Regents' Policy Manual.*