CS 362, HW 1

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Remember: you are encouraged to work on the homework in groups, but please observe the "Star Trek" rule from the syllabus.

1. Prove that $\log n! = \Theta(n \log n)$ and that $n! = \omega(2^n)$ and $n! = o(n^n)$

2. Assume you have functions f and g, such that $f(n)$ is $O(g(n))$. For each of the following statements, decide whether you think it is true or false and give either a proof or a counterexample

   (a) $\log_2 f(n)$ is $O(\log_2 g(n))$
   (b) $2^{f(n)}$ is $O(2^{g(n)})$
   (c) $f(n)^2$ is $O(g(n)^2)$

3. A c-tree is a tree with each node colored either red, green or silver that obeys the following rules:

   - Each red node has two children, exactly one of which is green.
   - Each green node has exactly one child, which is not green
   - Silver nodes have no children.

   Let R and S respectively denote the number of red and silver nodes, and n be the total number of nodes. Prove by induction that in any c-tree with $n \geq 1$, $S = R + 1$.

4. Write and solve a recurrence relation giving the number of strings of n digits containing at least one 3. For example, if $n = 5$, then 02309 would be one such string.
   In particular, let $f(n)$ be the number of strings of n digits with at least one 3. First, write an equation $f(n) = ***$, where the *** part contains smaller sub-problems, i.e. the $f(j)$ terms all have $j < n$. Then give a base case for the recurrence. Finally, use guess and check to solve the recurrence to within $\Theta()$ bounds. Hint: you may find the Master method useful for getting a good guess.