

## Conditional Distributions

A conditional distribution is the ratio of a joint distribution and a marginal distribution. When the value of random variable  $X$  is conditioned on the value of random variable  $Y$ :

$$p_{X|Y}(x | y) = \frac{p_{XY}(x, y)}{p_Y(y)}.$$

This can be generalized so that the values of  $N$  random variables  $X_1...X_N$  are conditioned on the values of  $M$  random variables  $Y_1...Y_M$ :

$$p_{X_1...X_N|Y_1...Y_M}(x_1...x_N | y_1...y_M) = \frac{p_{X_1...X_N, Y_1...Y_M}(x_1...x_N, y_1...y_M)}{p_{Y_1...Y_M}(y_1...y_M)}.$$

## Higher Order Markov Processes

Let  $S$  be a set of states:

$$S = \{1, 2, 3 \dots N\}$$

and let  $i, j, k \dots \in S$ . A random process is an order one Markov process iff:

$$p_{t|t-1,t-2\dots-\infty}(i|j,k\dots) = p_{t|t-1}(i|j).$$

The probability that the Markov process is in state  $i$  at time  $t$  is given by the following update formula:

$$p_t(i) = \sum_{j=1}^N p_{t|t-1}(i|j)p_{t-1}(j).$$

## Higher Order Markov Processes (contd.)

Let  $S$  be a set of states:

$$S = \{1, 2, 3 \dots N\}$$

and let  $i, j, k \dots \in S$ . A random process is an order two Markov process iff:

$$p_{t|t-1,t-2\dots-\infty}(i|j,k\dots) = p_{t|t-1,t-2}(i|j,k).$$

The probability that the Markov process is in state  $i$  at time  $t$  is given by the following update formula:

$$p_t(i) = \sum_{j=1}^N \sum_{k=1}^N p_{t|t-1,t-2}(i|j,k) p_{t-1,t-2}(j,k).$$

## Higher-Order Markov Processes (contd.)

A Markov process of order two can be thought of as a mapping between two joint distributions. Both of these joint distributions give the probability that the process visits two states in two successive times:

$$p_{t,t-1}(i, j) = \sum_{k=1}^N p_{t|t-1,t-2}(i | j, k) p_{t-1,t-2}(j, k).$$

The state  $i$  at time  $t$  is a marginal distribution (produced by summing over all possible states  $j$  at time  $t - 1$ ):

$$p_t(i) = \sum_{j=1}^N p_{t,t-1}(i, j).$$

## Higher Order Markov Processes (contd.)

It follows that a Markov process of order two, with states,  $S$ :

$$S = \{1, 2, 3 \dots N\}.$$

can be reduced to a Markov process of order one, with states,  $S' = S \times S$ :

$$S' = \{\langle 1, 1 \rangle, \langle 1, 2 \rangle \dots \langle N, N \rangle\}$$

and transition probability matrix:

$$p'_{t|t-1}(\langle i, j \rangle \mid \langle j, k \rangle) = p_{t|t-1,t-2}(i \mid j, k)$$

so that:

$$p'_t(\langle i, j \rangle) = \sum_{k=1}^N p'_{t|t-1}(\langle i, j \rangle \mid \langle j, k \rangle) p'_{t-1}(\langle j, k \rangle)$$

and

$$p_t(i) = \sum_{j=1}^N p'_t(\langle i, j \rangle).$$

## Information Source with Memory

An *information source with memory* generates messages using a source alphabet of length,  $M$ . If the source is modeled as a Markov process of order one, then the entropy of a message of length  $N$  is:

$$H_1 = H_0 + (N - 1)H_{t|t-1}$$

where

$$H_0 = - \sum_{i=1}^M p_t(i) \log p_t(i)$$

is the entropy of the first symbol and

$$H_{t|t-1} = - \sum_{i=1}^M \sum_{j=1}^M p_{t,t-1}(i, j) \log p_{t|t-1}(i | j)$$

is the entropy of each of the remaining symbols.

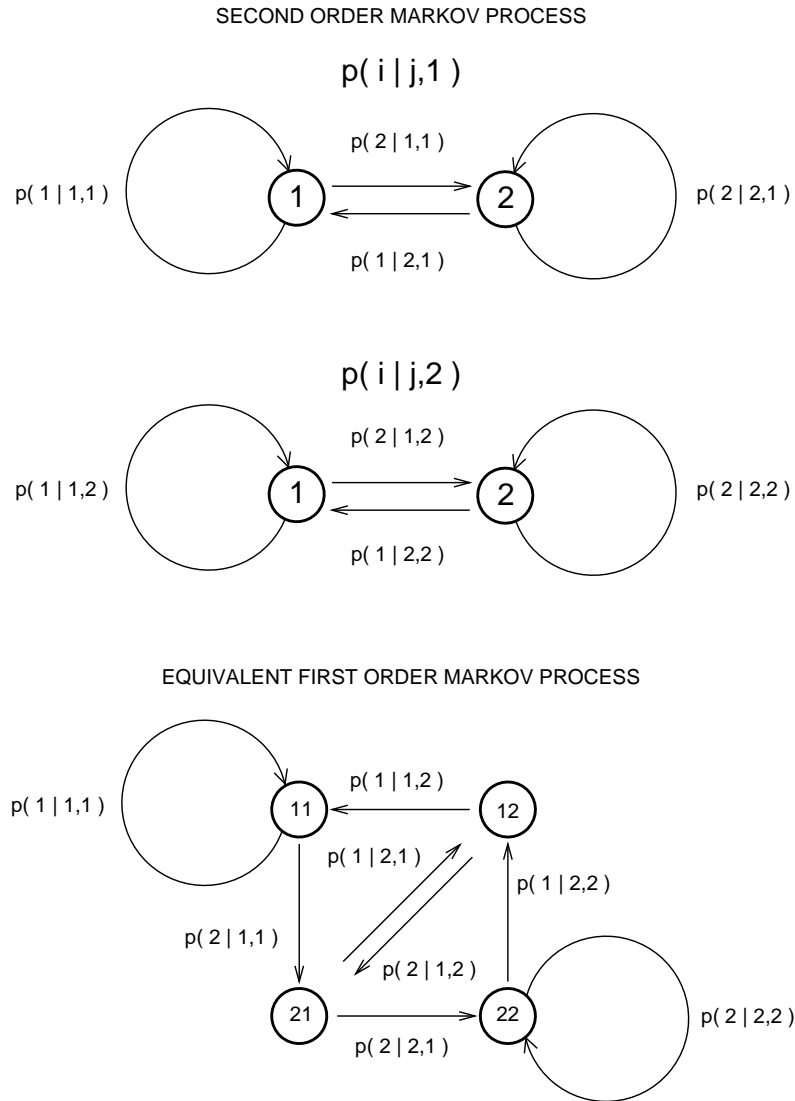


Figure 1: Second-order two-state Markov process and reduction to equivalent first-order Markov process.

## Example One

On avg., how much information is provided by each character in a random string of zeros and ones? The distribution for the  $t$ -th character is:

$$p_t(0) = 0.5$$

$$p_t(1) = 0.5$$

$$H_t = -0.5 \log(0.5) - 0.5 \log(0.5) = 1 \text{ bit.}$$

Each symbol delivers 1 bit of information on avg. in the memoryless case.



## Example Two

Now let's consider a string where the first character is chosen at random, but the remaining characters follow a simple pattern:

$$0101\dots01 \text{ or } 1010\dots10$$

The distribution for the  $t$ -th character is:

$$p_t(0) = 0.5$$

$$p_t(1) = 0.5$$

$$H_t = -0.5 \log(0.5) - 0.5 \log(0.5) = 1 \text{ bit.}$$

The joint distribution is:

$$\begin{bmatrix} p_{t,t-1}(0,0) & p_{t,t-1}(0,1) \\ p_{t,t-1}(1,0) & p_{t,t-1}(1,1) \end{bmatrix} = \begin{bmatrix} 0 & 0.5 \\ 0.5 & 0 \end{bmatrix}$$

## Example Two (contd.)

The conditional distribution is:

$$p_{t|t-1}(i|j) = \frac{p_{t,t-1}(i,j)}{p_{t-1}(j)}$$

$$\begin{bmatrix} p_{t|t-1}(0|0) & p_{t|t-1}(0|1) \\ p_{t|t-1}(1|0) & p_{t|t-1}(1|1) \end{bmatrix} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

and the conditional entropy per character is:

$$\begin{aligned} H_{t|t-1} &= - \sum_{i=0}^1 \sum_{j=0}^1 p_{t,t-1}(i,j) \log p_{t|t-1}(i|j) \\ &= -0.5 \log(1.0) + 0.5 \log(1.0) \\ &= 0 \text{ bits.} \end{aligned}$$

This is less than in the memoryless case.

## Information Source with Memory (contd.)

If the source is modeled as a Markov process of order two, then the entropy of a message of length  $N$  is:

$$H_2 = H_0 + H_{t|t-1} + (N - 2)H_{t|t-1,t-2}$$

where  $H_0$  and  $H_{t|t-1}$  are the entropies of the first and second symbols and

$$H_{t|t-1,t-2} = - \sum_{i=1}^M \sum_{j=1}^M \sum_{k=1}^M p_{t,t-1,t-2}(i, j, k) \log p_{t|t-1,t-2}(i | j, k)$$

is the entropy of each the remaining symbols.

## Information Limit

Let  $H_0$  be the entropy computed under the assumption that an information source is memoryless, and let  $H_1$  be the entropy computed under the assumption that the source is a Markov process of order one, and  $H_2$  be the entropy computed under the assumption that the source is a Markov process of order two, etc. Then

$$H_0 \geq H_1 \geq H_2 \geq \dots \geq \lim_{k \rightarrow \infty} H_k.$$