Notes for 8/31/09

f = (g + h) - (i + j)f, g, h, i, j := \$s0, \$s1, \$s2, \$s3, \$s4

How could we write this in MIPS assembly?

We will write this as though f...j are global variables, not stack variables.

```
add $t0, $s1, $s2
add $t1, $s3, $s4
sub $s0, $t0, $t1
la $t7, f
sw $s0, 0($t7)
A[12] = h + A[8]
A:= $s3
h := $s2
```

How could we write this in MIPS assembly?

To allocate an int array of size 20 (assume ints are 32 bit) as though it were a global variable, we could write this line in the .data section:

A: .space 80

Then in our .text section...

```
lw $t0, 32($s3)
add $t0, $s2, $t0
sw $t0, 48($s3)
A[i]
A:=$s3
i := $s4
```

How could we write this in MIPS assembly?

Branches in MIPS

if (a != b) { ... } else { ... } a, b := \$t0, \$t1

How could we write this in MIPS assembly?

How could we write this in MIPS assembly?

```
beq $t0, $t1, LabelA // branch if not equal this time
. . .
b End // jump unconditionally to End
LabelA:
. . .
End:
if (a < b) {
....
} else {
....
}
a, b := $t0, $t1
```

How could we write this in MIPS assembly?

Without using a branch pseudo-instruction:

```
slt $t7, $t0, $t1 // set $t7 to 1 if $t0 < $t1
bne $t7, $zero, LabelA // branch if $t7 is not zero
. . .
b End // jump unconditionally to End
LabelA:
. . .
End:</pre>
```

Using a branch pseudo-instruction:

```
blt $t0, $t1, LabelA // branch if $t0 < $t1
. . .
b End // jump unconditionally to End
LabelA:
. . .
End:
if (a <= b) {
...
} else {
...
}
a, b := $t0, $t1</pre>
```

How could we write this in MIPS assembly?

Without using a branch pseudo-instruction:

```
slt $t7, $t1, $t0 // set $t7 to 1 if $t1 < $t0
beq $t7, $zero, LabelA // branch if $t7 is not zero
. . .
b End // jump unconditionally to End
LabelA:
. . .
End:
```

Using a branch pseudo-instruction:

```
bgt $t1, $t0, LabelA // branch if $t1 > $t0
. . .
b End // jump unconditionally to End
LabelA:
. . .
End:
```