

A Complete Analysis of Resultants and Extraneous Factors for Unmixed Bivariate Polynomial Systems using the Dixon formulation *

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Abstract

A necessary and sufficient condition on the support of a generic unmixed bivariate polynomial system is identified such that for polynomial systems with such support, the Dixon resultant formulation produces their resultants. It is shown that Sylvester-type matrices can also be obtained for such polynomial systems. These results are shown to be a generalization of related results recently reported by Chionh as well as Zhang and Goldman. For a support not satisfying the above condition, the degree of the extraneous factor in the projection operator computed by the Dixon formulation is calculated by analyzing how much the support deviates from a related rectangular support satisfying the condition. The concept of the interior of a support is introduced; a generic inclusion of terms corresponding to support interior points in a polynomial system does not affect the nature of the projection operator computed by the Dixon construction.

For generic mixed bivariate systems, "good" Sylvester type matrices can be constructed by solving an optimization problem on their supports. The determinant of such a matrix gives a projection operator with a low degree extraneous factor. The results are illustrated on a variety of examples.

1 Introduction

New results characterizing generic unmixed polynomial systems with two variables for which resultants can be exactly computed and Sylvester-type matrices can be constructed are proved. Earlier in [CK00a], Chtcherba and Kapur had shown that the support of a bivariate unmixed polynomial system not including an orderable simplex is a necessary and sufficient condition for the determinant of the associated Dixon matrix being exact resultant (without any extraneous factors). Independently, Chionh [Chi01], Zhang and Goldman [ZG00], as well as Zhang in his Ph.D. thesis [Zha00] proposed corner-cut supports for which Dixon matrices as well as Sylvester-type multiplier matrices can be constructed whose determinant is the exact resultant. The results in this paper are shown to be related to and more general than those in [Chi01, ZG00, Zha00]. A necessary and sufficient condition on bivariate supports is identified such that for a generic unmixed polynomial system with such a support, its resultant can be computed exactly using constructions based on the Dixon resultant formulation. These bivariate supports are shown to include Chionh as well as Zhang and Goldman's corner-cut supports. In addition, for bivariate polynomial systems whose support do not satisfy the proposed conditions, the proposed construction also provides insight into the degree of the extraneous factor in a projection operator computed from a Dixon multiplier matrix.

The algorithm for constructing multiplier matrices based on the Dixon resultant formulation works in general for polynomial system from which more than two variables need to be eliminated even when the polynomial

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system is not necessarily unmixed. These multiplier matrices can be used to extract (in most cases) the resultants as determinants of their maximal minors. The approach generalizes a related method for constructing multiplier matrices from the Dixon resultant formulation discussed in a paper by Chtcherba and Kapur in RWCA'00. Beside being a generalization, the approach has the advantage of generating Sylvester-like multiplier matrices whose determinants are resultants even in cases where the earlier method by Chtcherba and Kapur produces an extraneous factor.

It is also shown that for the bivariate case, the proposed construction produces multiplier matrices with resultants as their determinants even in some mixed systems. The term in the support of a polynomial system is selected by formulating an optimization problem that minimizes the size of the Dixon multiplier matrix. The approach is compared with other approaches, and is shown to be more efficient and to work better on many examples of practical interest.

For unmixed polynomial systems in which more than two variables are simultaneously eliminated, the determinant of the associated Dixon multiplier matrix is shown to be not necessarily the resultant even for corner-cut supports. Nevertheless preliminary results show that with proper generalization to polynomial systems with more variables, the results will hold when more than two variables are eliminated.

Section 2 defines supports of a polynomial and a polynomial system, and reviews the BKK bound for toric roots and toric resultants of a polynomial system. Section 3 reviews the Dixon formulation of resultants, where the Dixon matrix and the Dixon polynomial of a given polynomial system are introduced. Section 4 analyzes the support of the Dixon polynomial in terms of the support of the polynomial system. It is shown that the support of the Dixon polynomial can be expressed as a union of the support of the Dixon polynomials of polynomial systems corresponding to simplexes. The concept of complement support characterizing how different a given support is from a bi-degree support (in the case of bivariate systems) is introduced; this support complement can be partitioned into four corners. It is shown that for a given polynomial system, the support of its Dixon polynomial (which determines the size of the Dixon matrix) can be shown to be a rectangle from which the four corner support complement determined from the support of the polynomial system are removed. In that section, points inside the convex hull of the support of a polynomial system are classified into two categories: (i) *support interior* points such that when terms corresponding to these support points are included in the polynomial system, the support of the Dixon polynomial and hence, the size of the Dixon matrix does not change, (ii) other support points such that the corresponding terms when included in the polynomial system contribute to the extraneous factors in the projection operator computed from the associated Dixon matrix.

Section 5 includes the main results of the paper, where the size of the Dixon matrix is related to the BKK bound. It is shown that the support of the polynomial system, after inclusion of its support interior points, is a rectangle with four rectangular corners removed, then the size of the Dixon matrix is the same as the BKK bound; this implies that for generic unmixed polynomial systems with such supports, the Dixon formulation computes the resultant exactly. For generic unmixed polynomial systems whose supports do not satisfy this condition, it is shown how the degree of the extraneous factor can be precisely estimated. In contrast, Chionh [Chi01] proved that the Dixon formulation computes the exact resultant for generic unmixed polynomial systems whose support is a rectangle with four rectangular corners removed. In Section 6, these results are extended to Dixon multiplier matrices, which are Sylvester type matrices but constructed using the Dixon formulation. Zhang and Goldman's results about corner cut supports are shown to be a special case of our results discussed in Sections 5 and 6. Section 7 considers mixed polynomial systems. It is first shown how some of the results from Sections 5 and 6 apply for a subclass of generic mixed polynomial systems as well. To generate "good" Dixon multiplier matrices whose determinants are projection operators with minimal degree extraneous factor, an optimization problem minimizing the support of the Dixon polynomial is formulated in terms of the term used to generate the Dixon multiplier matrix. An example illustrating this idea is discussed in detail. Section 8 compares our results experimentally with other approaches on examples.

2 Bivariate Systems

Consider a bivariate polynomial system \mathcal{F} ,

$$f_0 = \sum_{\alpha \in \mathcal{A}_0} a_{i,j} x^{\alpha_x} y^{\alpha_y}, \quad f_1 = \sum_{\beta \in \mathcal{A}_1} b_{i,j} x^{\beta_x} y^{\beta_y}, \quad f_2 = \sum_{\gamma \in \mathcal{A}_2} c_{i,j} x^{\gamma_x} y^{\gamma_y},$$

where for $i = 0, 1, 2$, each finite set \mathcal{A}_i of nonnegative integer tuples is called the *support* of the polynomial f_i ; it is denoted by $\mathcal{A}_i = S_x(f_i)$. If $\mathcal{A}_0 = \mathcal{A}_1 = \mathcal{A}_2$, the polynomial system is called *unmixed*; otherwise, it is called *mixed*.

We will denote the support of a polynomial system \mathcal{F} by $\mathcal{A} = \langle \mathcal{A}_0, \mathcal{A}_1, \mathcal{A}_2 \rangle$. Given a support \mathcal{A}_i , let $\text{Vol}(\mathcal{A}_i)$ stand for the Euclidean volume of the convex hull of \mathcal{A}_i .

Theorem 2.1 (BKK) *Given two bivariate polynomials f_1, f_2 , with corresponding supports \mathcal{A}_1 and \mathcal{A}_2 , the number of common toric roots of these polynomials is either infinite or at most*

$$\mu(\mathcal{A}_1, \mathcal{A}_2) = \text{Vol}(\mathcal{A}_1 + {}^1\mathcal{A}_2) - \text{Vol}(\mathcal{A}_1) - \text{Vol}(\mathcal{A}_2);$$

further, for most choices of coefficients, this bound is exact. The function μ is called the mixed volume function [GKZ94].

If $\mathcal{A}_1 = \mathcal{A}_2$, then $\mu(\mathcal{A}_1, \mathcal{A}_2) = 2\text{Vol}(\mathcal{A}_1)$.

In general, a polynomial system is called *generic* if it has a finite number of roots which is maximal for any choice of coefficients. The polynomial system $\{f_1, f_2\}$ is thus generic if the number of toric roots of any two polynomials equals its BKK bound. If we assume that coefficients are algebraically independent, then the polynomial system is certainly generic. Henceforth, the coefficients of terms in a polynomial system are assumed to be algebraically independent, unless stated otherwise.

In a generic case, the toric resultant of $\mathcal{F} = \{f_0, f_1, f_2\}$ is of degree equal to the BKK bound in terms of the coefficients of remaining polynomial f_0 [PS93]. For example, the degree of the resultant in terms of coefficients of f_0 is $\mu(\mathcal{A}_1, \mathcal{A}_2)$.

Using the Sylvester dialytic method, one can construct the resultant matrix for a given polynomial system by multiplying each polynomial by a set of monomials, called its *multipliers*, and rewriting the resulting polynomial system in the matrix notation. Let $X_i = \{x^a y^b\}$, $i = 0, 1, 2$, be the multiplier set for the polynomial f_i , respectively; then the matrix is constructed as

$$\begin{pmatrix} X_0 f_0 \\ X_1 f_1 \\ X_2 f_2 \end{pmatrix} = M \times X,$$

where X is the ordered set of all monomials appearing in $X_i f_i$ for $i = 0, 1, 2$. Note in order for M to qualify as a resultant matrix, $|X_0| \geq \mu_0 = \mu(\mathcal{A}_1, \mathcal{A}_2)$, $|X_1| \geq \mu_1 = \mu(\mathcal{A}_0, \mathcal{A}_2)$, and $|X_2| \geq \mu_2 = \mu(\mathcal{A}_0, \mathcal{A}_1)$.

If it can be shown that the matrix M above is square and non-singular, then it is the resultant matrix since the determinant of M is a multiple of the resultant. Moreover, if $|X_i| = \mu_i$, then M is exact, in the sense that its determinant is exactly the resultant of $\mathcal{F} = \{f_0, f_1, f_2\}$.

3 Dixon Resultant Matrix

In this section, we briefly review the generalized Dixon formulation, first introduced by Dixon [Dix08], and generalized by Kapur, Saxena and Yang [KSY94, KS96]. We will consider the bivariate case only.

¹The sum $\mathcal{A}_1 + \mathcal{A}_2$ is the Minkowski sum of polytopes $\mathcal{A}_1, \mathcal{A}_2$, where $p \in \mathcal{A}_1 + \mathcal{A}_2$ if $p = q + r$ for $q \in \mathcal{A}_1$ and $r \in \mathcal{A}_2$ where $+$ is the regular vector addition, see [CLO98] for definitions.

Define the *Dixon polynomial* to be

$$\theta_{x,y}(f_0, f_1, f_2) = \frac{1}{(\bar{x} - x)(\bar{y} - y)} \begin{vmatrix} f_0(x, y) & f_1(x, y) & f_2(x, y) \\ f_0(\bar{x}, y) & f_1(\bar{x}, y) & f_2(\bar{x}, y) \\ f_0(\bar{x}, \bar{y}) & f_1(\bar{x}, \bar{y}) & f_2(\bar{x}, \bar{y}) \end{vmatrix},$$

where \bar{x} and \bar{y} are new variables. Let \bar{X} be an ordered set of all monomials appearing in $\theta(f_0, f_1, f_2)$ in terms of variables \bar{x}, \bar{y} , and X be the set of all monomial in terms of variables x and y . Then

$$\theta_{x,y}(f_0, f_1, f_2) = \bar{X} \Theta_{x,y} X,$$

where $\Theta_{x,y}$ is called the *Dixon matrix*. Note that $\Theta_{x,y} = \Theta_{y,x}^T$, where the order of variables x, y is reversed; we will thus drop variable subscripts since it suffices to consider any variable order.

If $\mathcal{F} = \{f_0, f_1, f_2\}$ has a common zero, it is also a zero of $\theta(f_0, f_1, f_2)$ for any value of new variables \bar{x} and \bar{y} . Thus,

$$\Theta \times X = 0, \tag{1}$$

whenever x, y are replaced by a common zero of f_0, f_1, f_2 .

For polynomials $\{f_0, f_1, f_2\}$ to have a common zero, the equation (1) must be satisfied. If Θ is square and nonsingular, then its determinant must vanish, implying that under certain conditions, Θ is a resultant matrix. Even though this matrix is quite different from matrices constructed using the Sylvester dialytic method, there is a direct connection between the two which will be discussed later (see also [CK00b]).

We are interested in identifying conditions when the resultant matrix Θ is **exact**, i.e., its determinant is exactly (up to a constant factor) the resultant. Also, when it is not, we are interested in predicting the extraneous factor in the determinant of Θ (at the very least, the degree of the extraneous factor). We are thus interested in analyzing the size and structure of the monomial set X ; its size tells the number of columns in Θ and hence, whether or not, Θ is exact.

4 Support Structure of Dixon Polynomial

As $\mathcal{A} = \langle \mathcal{A}_0, \mathcal{A}_1, \mathcal{A}_2 \rangle$, by $\sigma \in \mathcal{A}$, we mean the triple $\langle \alpha, \beta, \gamma \rangle$ such that $\alpha \in \mathcal{A}_0$, $\beta \in \mathcal{A}_1$ and $\gamma \in \mathcal{A}_2$.

The Dixon polynomial above can be expressed using the Cauchy-Binet formula as a sum of Dixon matrices of 3-point set supports as shown below.

$$\theta(f_0, f_1, f_2) = \sum_{\sigma \in \mathcal{A}} \sigma(\mathbf{c}) \sigma(\mathbf{x}), \tag{2}$$

where

$$\sigma(\mathbf{c}) = \begin{vmatrix} a_\alpha & a_\beta & a_\gamma \\ b_\alpha & b_\beta & b_\gamma \\ c_\alpha & c_\beta & c_\gamma \end{vmatrix} \quad \text{and} \quad \sigma(\mathbf{x}) = \frac{1}{(\bar{x} - x)(\bar{y} - y)} \begin{vmatrix} x^{\alpha_x} y^{\alpha_y} & \bar{x}^{\alpha_x} y^{\alpha_y} & \bar{x}^{\alpha_x} \bar{y}^{\alpha_y} \\ x^{\beta_x} y^{\beta_y} & \bar{x}^{\beta_x} y^{\beta_y} & \bar{x}^{\beta_x} \bar{y}^{\beta_y} \\ x^{\gamma_x} y^{\gamma_y} & \bar{x}^{\gamma_x} y^{\gamma_y} & \bar{x}^{\gamma_x} \bar{y}^{\gamma_y} \end{vmatrix}.$$

In a generic case, where $\sigma(\mathbf{c}) \neq 0$, the support of the Dixon polynomial is the union of supports of $\sigma(\mathbf{x})$, where $\sigma(\mathbf{x})$ is the Dixon polynomial of the monomials corresponding to $\sigma = \langle \alpha, \beta, \gamma \rangle$.

As seen from the above formula, in the generic case, the support of the Dixon polynomial as well as the size of the Dixon matrix is completely determined by the support of the polynomial system \mathcal{F} ; we thus write the support of the Dixon polynomial as

$$\Delta_{\mathcal{A}} = S_{\mathbf{x}}(\theta(f_0, f_1, f_2)) \quad \text{where} \quad \mathcal{A}_0 = S_{\mathbf{x}}(f_0), \mathcal{A}_1 = S_{\mathbf{x}}(f_1), \mathcal{A}_2 = S_{\mathbf{x}}(f_2), \mathcal{A} = \langle \mathcal{A}_0, \mathcal{A}_1, \mathcal{A}_2 \rangle.$$

We will also denote by $\theta_{\mathcal{A}}$, the Dixon polynomial of the generic polynomial system with support \mathcal{A} .

Proposition 4.1 Given a generic polynomial system with support \mathcal{A} , $\Delta_{\mathcal{A}} = \bigcup_{\sigma \in \mathcal{A}} \Delta_{\sigma}$.

The above proposition is proved in its full generality for polynomial systems in many variables in [CK02b].

If $|\mathcal{A}_i| = 1$ for $i = \{0, 1, 2\}$, we can give a precise description of $\Delta_{\mathcal{A}}$: it is the union of two disjoint adjacent rectangles as discussed below.

Definition 4.1 A box support \mathcal{B}_u^v , where $u, v \in \mathbb{N}^2$, is a dense rectangular support located at u , and is defined as

$$\mathcal{B}_u^v = \{ (i, j) \mid u_x \leq i < v_x \text{ and } u_y \leq j < v_y \}.$$

Given a support $\mathcal{A} = \langle \mathcal{A}_0, \mathcal{A}_1, \mathcal{A}_2 \rangle$, let $\mathcal{B}_{\mathcal{A}}$ stand for the smallest box support containing all \mathcal{A}_i 's.

Proposition 4.2 Let $\sigma = \{\alpha, \beta, \gamma\}$ be a 3-point support such that $\alpha_x \leq \beta_x \leq \gamma_x$ (w.l.o.g.); consider a generic polynomial system with the support σ . Then

$$\Delta_{\sigma} = \Delta_L \cup \Delta_H \quad \text{and} \quad \Delta_L \cap \Delta_H = \emptyset \quad \text{where} \quad \Delta_L = \mathcal{B}_s^t \quad \text{and} \quad \Delta_H = \mathcal{B}_p^q,$$

where $L = \{\alpha, \beta, (\beta_x, \gamma_y)\}$, $H = \{(\beta_x, \alpha_y), \beta, \gamma\}$, and

$$\begin{aligned} s_x &= \alpha_x, & p_x &= \beta_x, \\ t_x &= \beta_x, & q_x &= \gamma_x, \\ s_y &= \alpha_y + \min(\beta_y, \gamma_y), & p_y &= \gamma_y + \min(\alpha_y, \beta_y), \\ t_y &= \alpha_y + \max(\beta_y, \gamma_y), & q_y &= \gamma_y + \max(\alpha_y, \beta_y). \end{aligned}$$

Proof: The Dixon polynomial θ_{σ} of σ has the same monomials even when $\bar{x} = 1$ and $\bar{y} = 1$. Substituting $\bar{x} = 1$ and $\bar{y} = 1$ thus does not eliminate any monomials in x, y due to cancellation. So, $S_{\mathbf{x}}(\theta_{\sigma}) = S_{\mathbf{x}}(\theta_{\sigma}|_{\bar{x}=1, \bar{y}=1})$. Since

$$\begin{aligned} \begin{vmatrix} x^{\alpha_x} y^{\alpha_y} & y^{\alpha_y} & 1 \\ x^{\beta_x} y^{\beta_y} & y^{\beta_y} & 1 \\ x^{\gamma_x} y^{\gamma_y} & y^{\gamma_y} & 1 \end{vmatrix} &= \left(\begin{vmatrix} x^{\alpha_x} y^{\alpha_y} & y^{\alpha_y} & 1 \\ x^{\beta_x} y^{\beta_y} & y^{\beta_y} & 1 \\ x^{\beta_x} y^{\gamma_y} & y^{\gamma_y} & 1 \end{vmatrix} + \begin{vmatrix} x^{\beta_x} y^{\alpha_y} & y^{\alpha_y} & 1 \\ x^{\beta_x} y^{\beta_y} & y^{\beta_y} & 1 \\ x^{\gamma_x} y^{\gamma_y} & y^{\gamma_y} & 1 \end{vmatrix} \right) \\ &= (x^{\alpha_x} - x^{\beta_x})(y^{\alpha_y + \beta_y} - y^{\alpha_y + \gamma_y}) + (x^{\beta_x} - x^{\gamma_x})(y^{\beta_y + \gamma_y} - y^{\alpha_y + \gamma_y}), \end{aligned}$$

$$S_{\mathbf{x}}(\theta_{\sigma}) = S_{\mathbf{x}}\left(\frac{(x^{\alpha_x} - x^{\beta_x}) y^{\alpha_y} (y^{\beta_y} - y^{\gamma_y})}{1-x} + \frac{(x^{\beta_x} - x^{\gamma_x}) y^{\gamma_y} (y^{\beta_y} - y^{\alpha_y})}{1-x}\right).$$

The monomials coming from the two fractions are disjoint. Further, the support of each fraction is \mathcal{B}_s^t and \mathcal{B}_p^q , respectively. \square .

4.1 Support of Dixon Polynomial: A Generic System with Box Support

An unmixed polynomial system such that $\mathcal{A}_i = \mathcal{B}_u^v$ for some $u, v \in \mathbb{N}$ and $u = (0, 0)$, is called a **bi-degree** system. We allow u to be arbitrary, as one can shift the entire support to the origin, compute the Dixon polynomial of the new *shifted* polynomial system, and then multiply the result with the appropriate monomial.

Dixon in [Dix08] generalized the Bezout method for dense bi-degree polynomial systems, and established that matrices constructed using that method are exact, i.e., their determinants are the resultants of the polynomial systems. The support structure of the Dixon polynomial has been known, and is generalized to the n -degree systems in [KS96] and [Sax97]. Here, we give a precise description of a bi-degree system.

Proposition 4.3 The support of the Dixon polynomial of a bi-degree system is a rectangle

$$\Delta_{\mathcal{B}_u^v} = \mathcal{B}_s^t \quad \text{where} \quad t = (v_x - 1, 2v_y - 1) \quad \text{and} \quad s = (u_x, 2u_y).$$

Proof: $\Delta_{\mathcal{B}_u^v} \subseteq \mathcal{B}_u^t$ as the exponents of monomials of $\theta_{\mathcal{B}_u^v}$ cannot be outside \mathcal{B}_u^t . The expansion of the Dixon polynomial $\theta_{\mathcal{B}_u^v}$ has two simplexes among others, which include all support points in \mathcal{B}_u^t . Let $\sigma_1 = \langle u, v, (v_x, u_y) \rangle$ and $\sigma_2 = \langle u, v, (u_x, v_y) \rangle$. Using Proposition 4.2, it can be seen that these simplexes result in the entire \mathcal{B}_u^t . \square

The size of $\Delta_{\mathcal{B}_u^v}$ of the box support \mathcal{B}_u^v is $2(v_x - u_x)(v_y - u_y)$, which is precisely the mixed volume of any two polynomials with \mathcal{B}_u^v as support; hence, the Dixon matrix is exact for such a polynomial system.

An important point about box supports is that points “inside” the box support do not play any role in determining the support of the Dixon polynomial (which can be seen, from the proof of the above proposition). Later, we will give a precise description of points which do not influence the support of the Dixon polynomial. Identifying such “interior” points and not using them in computations can reduce the cost of algorithms based on this method.

4.2 Support Complement

Given an arbitrary unmixed support \mathcal{A} , we can analyze how different it is from a bi-degree system having $\mathcal{B}_\mathcal{A}$ as its support. We will define complement of \mathcal{A} w.r.t. $\mathcal{B}_\mathcal{A}$ in terms of simplexes making up \mathcal{A} .

Definition 4.2 *Given an unmixed polynomial system with \mathcal{A} as its support, the support complement of \mathcal{A} w.r.t. $\mathcal{B}_\mathcal{A}$ is defined as the set $\mathcal{C}_\mathcal{A}$ that can be expressed as:*

$$\mathcal{C}_\mathcal{A} = \mathcal{C}_{00}^{\mathcal{A}} \cup \mathcal{C}_{01}^{\mathcal{A}} \cup \mathcal{C}_{10}^{\mathcal{A}} \cup \mathcal{C}_{11}^{\mathcal{A}},$$

where each $\mathcal{C}_{i,j}^{\mathcal{A}}$, $i, j \in \{0, 1\}$, is called a corner complement-support of \mathcal{A} , and each $\mathcal{C}_{i,j}^{\mathcal{A}}$ satisfies the property that if $c \in \mathcal{C}_{i,j}^{\mathcal{A}}$, then there does not exist $p \in \mathcal{A}$ such that

$$\begin{aligned} p_x \leq c_x \text{ if } i = 0 & \quad \text{and} \quad p_x \geq c_x \text{ if } i = 1, \\ p_y \leq c_y \text{ if } j = 0 & \quad \text{and} \quad p_y \geq c_y \text{ if } j = 1. \end{aligned}$$

The complement of a support can also be defined as the intersection of simplexes making up the support, which with some technicalities, can be extended to mixed supports as well. The next proposition establishes the equivalence of the two definitions for the unmixed case.

Proposition 4.4 *The support complement $\mathcal{C}_\mathcal{A}$ of \mathcal{A} w.r.t. $\mathcal{B}_\mathcal{A}$, where $\mathcal{C}_\mathcal{A} = \mathcal{C}_{00}^{\mathcal{A}} \cup \mathcal{C}_{01}^{\mathcal{A}} \cup \mathcal{C}_{10}^{\mathcal{A}} \cup \mathcal{C}_{11}^{\mathcal{A}}$ and $0 \leq i, j \leq 1$,*

$$\mathcal{C}_{i,j}^{\mathcal{A}} = \bigcap_{\sigma \in \mathcal{A}} \mathcal{C}_{i,j}^{\sigma}.$$

Proof: Any point $c \in \bigcap_{\sigma \in \mathcal{A}} \mathcal{C}_{i,j}^{\sigma}$ satisfies the definition for being in a corner complement-support of \mathcal{A} . If $c \in \mathcal{C}_{i,j}^{\mathcal{A}}$, there does not exist $p \in \mathcal{A}$ preventing c from being in $\mathcal{C}_{i,j}^{\mathcal{A}}$; hence no simplex $\sigma \in \mathcal{A}$ can include p and therefore, $c \in \mathcal{C}_{i,j}^{\sigma}$ for some $\sigma \in \mathcal{A}$. \square

Figure 1 shows the support complement for an unmixed support \mathcal{A} , where the four corners are outlined. As the reader would notice, $\mathcal{C}_{10}^{\mathcal{A}}$ and $\mathcal{C}_{01}^{\mathcal{A}}$ overlap in general.

Definition 4.3 *Given an arbitrary box support \mathcal{B}_u^v and a support complement $\mathcal{C}_\mathcal{A} = \mathcal{C}_{00}^{\mathcal{A}} \cup \mathcal{C}_{01}^{\mathcal{A}} \cup \mathcal{C}_{10}^{\mathcal{A}} \cup \mathcal{C}_{11}^{\mathcal{A}}$ of some support \mathcal{A} , let the support $\mathcal{B}_u^v / \langle \mathcal{C}_{00}^{\mathcal{A}}, \mathcal{C}_{01}^{\mathcal{A}}, \mathcal{C}_{10}^{\mathcal{A}}, \mathcal{C}_{11}^{\mathcal{A}} \rangle$ stand for the box support with four corner complement-supports $\langle \mathcal{C}_{00}^{\mathcal{A}}, \mathcal{C}_{01}^{\mathcal{A}}, \mathcal{C}_{10}^{\mathcal{A}}, \mathcal{C}_{11}^{\mathcal{A}} \rangle$ removed from it. More precisely, the points $(0, 0)$ of $\mathcal{C}_{00}^{\mathcal{A}}$, $(0, \max_y)$ of $\mathcal{C}_{01}^{\mathcal{A}}$, $(\max_x, 0)$ of $\mathcal{C}_{10}^{\mathcal{A}}$, and (\max_x, \max_y) of $\mathcal{C}_{11}^{\mathcal{A}}$ are located, respectively, at (u_x, u_y) , (u_x, v_y) , (v_x, u_y) , (v_x, v_y) of \mathcal{B}_u^v .*

Using these notions, a connection between the support \mathcal{A} of a polynomial system and the support $\Delta_\mathcal{A}$ of its Dixon polynomial can be established.

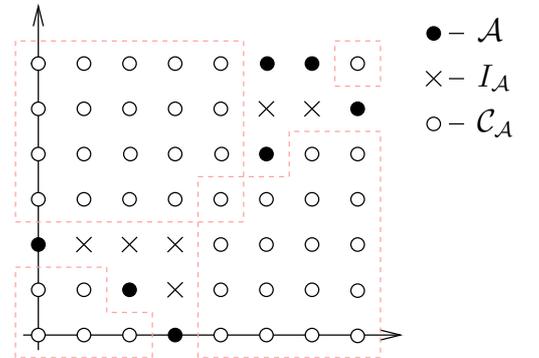


Figure 1: Complement Supports

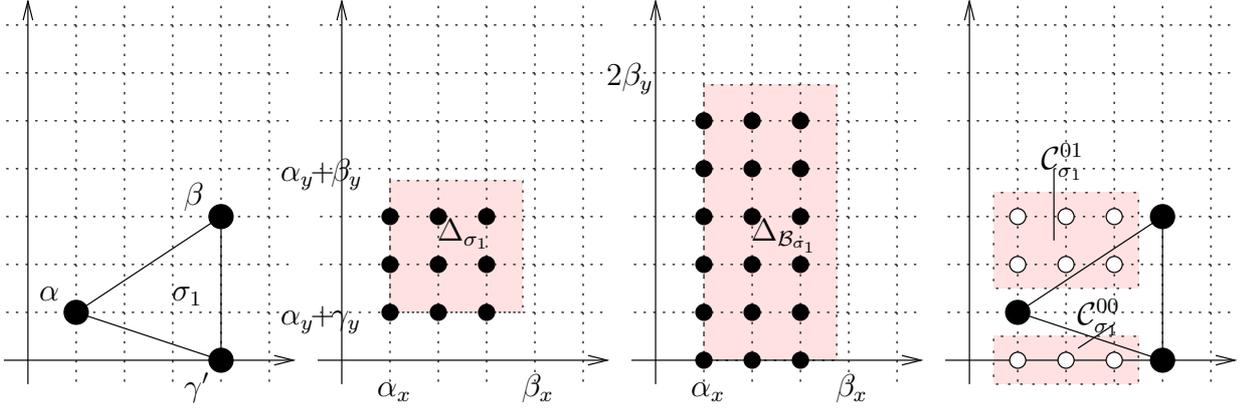


Figure 2: Example of σ_1 , Δ_{σ_1} , $\Delta_{\mathcal{B}_{\sigma_1}}$ and \mathcal{C}_{σ_1} used in the proof of Proposition 4.5.

4.3 Support of Dixon Polynomial: General Case

An interesting property of the support of the Dixon polynomial is that it admits a concise geometric description, given that it is a union of the supports of the Dixon polynomial of polynomial systems with smaller support sets. The following theorem gives the support of the Dixon polynomial of a bivariate polynomial system in terms of how different the support of the polynomial system is from the bi-degree support. It also enables computing the support of the Dixon polynomial without expanding all determinants in the formula for the Dixon polynomial. Figure 3 shows the support of the Dixon polynomial for the unmixed polynomial system shown in Figure 1.

Proposition 4.5 *Given a simplex σ , $\Delta_{\sigma} = \Delta_{\mathcal{B}_{\sigma}} / \langle \mathcal{C}_{00}^{\sigma}, \mathcal{C}_{01}^{\sigma}, \mathcal{C}_{10}^{\sigma}, \mathcal{C}_{11}^{\sigma} \rangle$.*

Proof: Let $\sigma = \langle \alpha, \beta, \gamma \rangle$ be a simplex; w.l.o.g., assume $\alpha_x \leq \beta_x \leq \gamma_x$. Let $\alpha' = (\beta_x, \alpha_y)$ and $\gamma' = (\beta_x, \gamma_y)$, where $\sigma_1 = \langle \alpha, \beta, \gamma' \rangle$ and $\sigma_2 = \langle \alpha', \beta, \gamma \rangle$. By Proposition 4.2,

$$\Delta_{\sigma} = \Delta_{\sigma_1} \cup \Delta_{\sigma_2} \quad \text{and} \quad \Delta_{\sigma_1} \cap \Delta_{\sigma_2} = \emptyset.$$

Further, $\Delta_{\mathcal{B}_{\sigma}} = \Delta_{\mathcal{B}_{\sigma_1}} \cup \Delta_{\mathcal{B}_{\sigma_2}}$ and $\Delta_{\mathcal{B}_{\sigma_1}} \cap \Delta_{\mathcal{B}_{\sigma_2}} = \emptyset$. The support complements of σ_1 and σ_2 are disjoint as well, and their union is the support complement of σ .

$$\mathcal{C}_{ij}^{\sigma} = \mathcal{C}_{ij}^{\sigma_1} \cup \mathcal{C}_{ij}^{\sigma_2} \quad \text{for all } i, j \in \{0, 1\}.$$

Therefore,

$$\Delta_{\mathcal{B}_{\sigma}} / \langle \mathcal{C}_{00}^{\sigma}, \mathcal{C}_{01}^{\sigma}, \mathcal{C}_{10}^{\sigma}, \mathcal{C}_{11}^{\sigma} \rangle = (\Delta_{\mathcal{B}_{\sigma_1}} / \langle \mathcal{C}_{00}^{\sigma_1}, \mathcal{C}_{01}^{\sigma_1}, \mathcal{C}_{10}^{\sigma_1}, \mathcal{C}_{11}^{\sigma_1} \rangle) \cup (\Delta_{\mathcal{B}_{\sigma_2}} / \langle \mathcal{C}_{00}^{\sigma_2}, \mathcal{C}_{01}^{\sigma_2}, \mathcal{C}_{10}^{\sigma_2}, \mathcal{C}_{11}^{\sigma_2} \rangle).$$

Hence we need to show that

$$\Delta_{\sigma_1} = \Delta_{\mathcal{B}_{\sigma_1}} / \langle \mathcal{C}_{00}^{\sigma_1}, \mathcal{C}_{01}^{\sigma_1}, \mathcal{C}_{10}^{\sigma_1}, \mathcal{C}_{11}^{\sigma_1} \rangle \quad \text{and} \quad \Delta_{\sigma_2} = \Delta_{\mathcal{B}_{\sigma_1}} / \langle \mathcal{C}_{00}^{\sigma_2}, \mathcal{C}_{01}^{\sigma_2}, \mathcal{C}_{10}^{\sigma_2}, \mathcal{C}_{11}^{\sigma_2} \rangle.$$

For σ_1 , there are 3 cases: $\alpha_y \leq \min(\beta_y, \gamma_y)$, $\alpha_y \geq \max(\beta_y, \gamma_y)$ or α_y is in between of β_y and γ_y . The last case is shown in Figure 2. Each case is similar, and here we only show the case when α_y is in between of β_y and γ_y . By Proposition 4.2, Δ_{σ_1} is a rectangular support, as well as $\Delta_{\mathcal{B}_{\sigma_1}}$ with boundaries shown in Figure 2. It can be established that either $\mathcal{C}_{ij}^{\sigma_1}$ is empty or is a rectangle having the width of $\beta_x - \alpha_x - 1$. In that case, it can be seen that $\mathcal{C}_{10}^{\sigma_1} = \mathcal{C}_{11}^{\sigma_1} = \emptyset$ and

$$\mathcal{C}_{00}^{\sigma_1} = \mathcal{B}_s^t \quad \text{where} \quad \begin{aligned} s &= (\alpha_x, \min(\beta_y, \gamma_y)), \\ t &= (\beta_x, \alpha_y), \end{aligned}$$

and

$$\mathcal{C}_{01}^{\sigma_1} = \mathcal{B}_u^v \quad \text{where} \quad \begin{aligned} u &= (\alpha_x, \alpha_y + 1), \\ v &= (\beta_x, \max(\beta_y, \gamma_y)). \end{aligned}$$

Shifting $\mathcal{C}_{00}^{\sigma_1}$ to $(\alpha_x, \min(\beta_y, \gamma_y))$ and shifting $\mathcal{C}_{01}^{\sigma_1}$ to $(\alpha_x, \max(\beta_y, \gamma_y) + \alpha_y + 1)$, and removing from $\Delta_{\mathcal{B}_{\sigma_1}}$ we get Δ_{σ_1} . Similarly, all other cases can be verified. \square

The following theorem generalizes Proposition 4.5.²

Theorem 4.1 *For any given unmixed polynomial system with the support \mathcal{A} and $\mathcal{C}_{\mathcal{A}} = \mathcal{C}_{00}^{\mathcal{A}} \cup \mathcal{C}_{01}^{\mathcal{A}} \cup \mathcal{C}_{10}^{\mathcal{A}} \cup \mathcal{C}_{11}^{\mathcal{A}}$ as the complement of \mathcal{A} with respect to $\mathcal{B}_{\mathcal{A}}$, $\Delta_{\mathcal{A}} = \Delta_{\mathcal{B}_{\mathcal{A}}} / \langle \mathcal{C}_{00}^{\mathcal{A}}, \mathcal{C}_{01}^{\mathcal{A}}, \mathcal{C}_{10}^{\mathcal{A}}, \mathcal{C}_{11}^{\mathcal{A}} \rangle$.*

Proof: We need to show that

$$\Delta_{\mathcal{B}_{\mathcal{A}}} / \langle \mathcal{C}_{00}^{\mathcal{A}}, \mathcal{C}_{01}^{\mathcal{A}}, \mathcal{C}_{10}^{\mathcal{A}}, \mathcal{C}_{11}^{\mathcal{A}} \rangle = \bigcup_{\sigma \in \mathcal{A}} \Delta_{\mathcal{B}_{\sigma}} / \langle \mathcal{C}_{00}^{\sigma}, \mathcal{C}_{01}^{\sigma}, \mathcal{C}_{10}^{\sigma}, \mathcal{C}_{11}^{\sigma} \rangle,$$

as using Proposition 4.5 and the fact that

$$\Delta_{\mathcal{A}} = \bigcup_{\sigma \in \mathcal{A}} \Delta_{\sigma},$$

the statement of the theorem follows. Let $Q_{\mathcal{A}} = \Delta_{\mathcal{B}_{\mathcal{A}}} / \langle \mathcal{C}_{00}^{\mathcal{A}}, \mathcal{C}_{01}^{\mathcal{A}}, \mathcal{C}_{10}^{\mathcal{A}}, \mathcal{C}_{11}^{\mathcal{A}} \rangle$ and correspondingly for $\sigma \in \mathcal{A}$, $Q_{\sigma} = \Delta_{\mathcal{B}_{\sigma}} / \langle \mathcal{C}_{00}^{\sigma}, \mathcal{C}_{01}^{\sigma}, \mathcal{C}_{10}^{\sigma}, \mathcal{C}_{11}^{\sigma} \rangle$. We need to show

$$Q_{\mathcal{A}} = \bigcup_{\sigma \in \mathcal{A}} Q_{\sigma}.$$

- Case $Q_{\sigma} \subseteq Q_{\mathcal{A}}$. Let $p \in Q_{\sigma}$ where $\sigma = \{\alpha, \beta, \gamma\} \in \mathcal{A}$. Clearly $p \in \Delta_{\mathcal{B}_{\sigma}}$, since $\Delta_{\mathcal{B}_{\sigma}} \subseteq \Delta_{\mathcal{B}_{\mathcal{A}}}$, we need to show that $p \notin \{\mathcal{C}_{00}^{\mathcal{A}}, \mathcal{C}_{01}^{\mathcal{A}}, \mathcal{C}_{10}^{\mathcal{A}}, \mathcal{C}_{11}^{\mathcal{A}}\}$ in which case $p \in Q_{\mathcal{A}}$. But this is clear since $p \notin \{\mathcal{C}_{00}^{\sigma}, \mathcal{C}_{01}^{\sigma}, \mathcal{C}_{10}^{\sigma}, \mathcal{C}_{11}^{\sigma}\}$, and $\mathcal{C}_{i,j}^{\mathcal{A}} \subseteq \mathcal{C}_{i,j}^{\sigma}$ in $\mathcal{B}_{\mathcal{A}}$.
- Case $Q_{\mathcal{A}} \subseteq \bigcup_{\sigma \in \mathcal{A}} Q_{\sigma}$. Let $p \in Q_{\mathcal{A}}$ we need to find $\sigma = \{\alpha, \beta, \gamma\} \in \mathcal{A}$ such that $p \in Q_{\sigma}$. Since $p \in Q_{\mathcal{A}}$, there must exist $\alpha, \beta, \gamma, \delta \in \mathcal{A}$ such that

$$\begin{aligned} \alpha_x &\leq p_x < \beta_x, & \text{and} \\ \gamma_y &\leq p_y < \gamma_y + \delta_y. \end{aligned} \quad (3)$$

Moreover, $p \in Q_{\{\alpha, \beta, \gamma, \delta\}}$. It should be clear that in order to satisfy equations of (3), only three points are necessary as depending on a configuration, either α can be replaced by γ or δ , or β can be replaced by γ or δ . \square

$\Delta_{\mathcal{A}}$ is dependent on the variable order used in $\theta_{\mathcal{A}}$, but the size of $\Delta_{\mathcal{A}}$ is the same for any variable order if \mathcal{A} is unmixed (by Theorem 4.1). The number of columns is determined by the size of the support in terms of variables x, y . On the other hand, the number of rows is determined by the size of support in terms of variables \bar{x}, \bar{y} , which is the same as if the variable order is reversed and the support is considered in terms of variables x, y .

The above observation thus implies that the Dixon matrix is square for unmixed polynomial systems, but it need not be square for mixed polynomial systems.

²We recently learned from Prof. Goldman that the same theorem appeared in [Chi01] and was presented in a different way. The proof used in [Chi01] is quite different from our proof.

4.4 Support Interior

In this subsection, the role of points in the convex hull of a given support \mathcal{A} vis a vis their contribution to the resultant and extraneous factors generated by the Dixon formulation are investigated. Points strictly inside in the convex hull of a support can be classified into two broad categories: (i) terms corresponding to those points whose inclusion into the polynomial system does not change the resultant (or the projection operator), and (ii) terms corresponding to those points whose inclusion into the polynomial system changes the projection operator and can result in more extraneous factors. First, the concept of an *interior point* relative to a support is introduced. Intuitively, an interior point is a point whose presence in the support of a polynomial in a given polynomial system does not change the support of its Dixon polynomial.

Definition 4.4 A point p is called **support-interior** w.r.t an unmixed support \mathcal{A} if and only if $\exists \{\alpha, \beta, \gamma, \delta\} \in \mathcal{A}_i$ for some i , such that

$$\begin{aligned} \alpha_x \leq p_x \leq \beta_x & \quad \text{and} \quad \gamma_x \leq p_x \leq \delta_x, \\ \alpha_y \leq p_y \leq \beta_y & \quad \text{and} \quad \gamma_y \geq p_y \geq \delta_y. \end{aligned}$$

Definition 4.4 is geometric that allows to verify a given point for being in support interior with respect to \mathcal{A} . A support interior point does not have to be in \mathcal{A} .

Given \mathcal{A} , let $I_{\mathcal{A}}$ consist of all support-interior points of \mathcal{A} and $\overline{\mathcal{A}} = \mathcal{A} \cup I_{\mathcal{A}}$; $I_{\mathcal{A}}$ is called the *support-interior* of \mathcal{A} . For example, all points of the box support \mathcal{B}_u^v are support-interior with the exception of the four corner points.

It is easy to see that a support-interior point is in the convex hull of \mathcal{A} , but the converse is not necessarily the case. For example, the convex hull of support $\{(0, 0), (0, 2), (2, 0)\}$ contains the point $(1, 1)$, but it is not a support-interior point. Support-interior of \mathcal{A} is also not the same as the convex hull interior, but is a subset of the convex hull interior as shown below.

Proposition 4.6 Given a support \mathcal{A} , its support interior $I_{\mathcal{A}}$ is contained in the Newton polytope of \mathcal{A} , i.e.

$$I_{\mathcal{A}} \subset \text{cHull}(\mathcal{A}) \cap \mathbb{N}^2.$$

Proof: Consider any support-interior point $p = (p_x, p_y)$ in \mathcal{A} . By definition, there are points a, b, c, d satisfying the above stated condition on the coordinates of p, a, b, c, d . Since (i) the intersection points of line $x = p_x$ with lines ad and bc are in the convex hull and (ii) p lies in between these intersection points, p is in the convex hull. \square .

In the construction of $\Delta_{\mathcal{A}}$, it does not matter whether \mathcal{A} is used or $\overline{\mathcal{A}}$ is used. In other words, extending \mathcal{A} with its support-interior points will not change $\Delta_{\mathcal{A}}$.

Theorem 4.2 Given a generic unmixed polynomial system with support \mathcal{A} ,

$$\Delta_{\mathcal{A}} = \Delta_{\overline{\mathcal{A}}}.$$

Proof: The presence of a support interior point does not change $\mathcal{B}_{\mathcal{A}}$, nor any of $\mathcal{C}_{i,j}^{\mathcal{A}}$; therefore, it has no effect on the support $\Delta_{\mathcal{A}}$ of the Dixon polynomial.

On the other hand, if a point is not support interior, then it belongs to one of $\mathcal{C}_{i,j}^{\mathcal{A}}$, and hence changes $\Delta_{\mathcal{A}}$. \square .

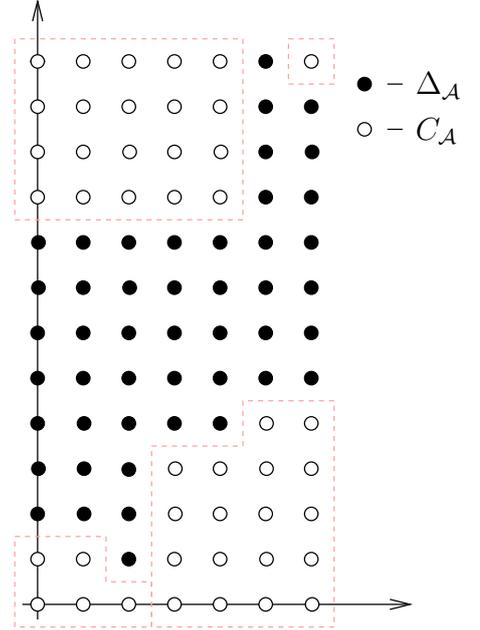


Figure 3: $\Delta_{\mathcal{A}}$ where \mathcal{A} is from Figure 1

5 Dixon Matrix Size: Exact Cases vs Extraneous Factors

In this section, we relate the size of the Dixon matrix associated with a given polynomial system, which is determined by the size of the support of its Dixon polynomial, to the BKK bound on the number of its toric roots, which is determined by the mixed volume of the Newton polytopes of its supports. We identify necessary and sufficient conditions on the support of the polynomial system under which the Dixon matrix is exact in the sense that its size is precisely the BKK bound. When these conditions on the support are not satisfied, we give an estimate on the degree of the extraneous factor in the projection operator extracted from the Dixon matrix by relating its size to the BKK bound.

Theorem 5.1 *The size of the Dixon matrix of an unmixed polynomial system $\mathcal{F} = \{f_0, f_1, f_2\}$ with support \mathcal{A} is*

$$|\Delta_{\mathcal{A}}| = |\Delta_{\mathcal{B}_{\mathcal{A}}}| - |\mathcal{C}_{00}^{\mathcal{A}}| - |\mathcal{C}_{01}^{\mathcal{A}}| - |\mathcal{C}_{10}^{\mathcal{A}}| - |\mathcal{C}_{11}^{\mathcal{A}}|.$$

Proof: Note that inside $\Delta_{\mathcal{B}_{\mathcal{A}}}$, $\mathcal{C}_{i,j}^{\mathcal{A}}$ do not intersect with each other for $i, j \in \{0, 1\}$, i.e., $\mathcal{C}_{00}^{\mathcal{A}}$ and $\mathcal{C}_{11}^{\mathcal{A}}$ do not intersect $\mathcal{C}_{10}^{\mathcal{A}}$ or $\mathcal{C}_{01}^{\mathcal{A}}$. Even if within $\mathcal{B}_{\mathcal{A}}$, $\mathcal{C}_{00}^{\mathcal{A}}$ intersects $\mathcal{C}_{11}^{\mathcal{A}}$, they cannot intersect within $\Delta_{\mathcal{B}_{\mathcal{A}}}$ as $\Delta_{\mathcal{B}_{\mathcal{A}}}$ is almost twice the size of $\mathcal{B}_{\mathcal{A}}$, and $\mathcal{C}_{00}^{\mathcal{A}}$ is obviously smaller than $\mathcal{B}_{\mathcal{A}}$. This is true for $\mathcal{C}_{01}^{\mathcal{A}}$ and $\mathcal{C}_{10}^{\mathcal{A}}$ as well. The statement of the theorem follows directly from Theorem 4.1 and the definition of $\Delta_{\mathcal{A}}$. \square .

The same theorem is independently proved in [Chi01]; the proof method seems to be quite different, however.

How does the size of the Dixon matrix compare with the BKK bound of \mathcal{F} ? For the unmixed case, if the size of the Dixon matrix equals the BKK bound of any two polynomials in \mathcal{F} , then the matrix is exact, i.e., its determinant is exactly the toric resultant. To see the relationship between the BKK bound which is defined in terms of Newton polytopes, and the size of the Dixon matrix, we can characterize how different the convex hull of the support \mathcal{A} is from the box support using *corners*. Let

$$\mathcal{Q} = \text{Conv}_{\text{hull}}(\mathcal{B}_{\mathcal{A}}) - \text{Conv}_{\text{hull}}(\mathcal{A}),$$

which can be split into four disjoint, not necessarily convex, polyhedral sets:

$$\mathcal{Q} = \mathcal{Q}_{00}^{\mathcal{A}} \cup \mathcal{Q}_{01}^{\mathcal{A}} \cup \mathcal{Q}_{10}^{\mathcal{A}} \cup \mathcal{Q}_{11}^{\mathcal{A}}.$$

Figure 4 shows the Newton polytope complement for the earlier example shown in Figure 3.

For an unmixed polynomial system, the BKK bound, which is the mixed volume of any two polynomials with the support \mathcal{A} , is

$$\mu(\mathcal{A}_0, \mathcal{A}_1) = 2\text{Vol}(\mathcal{A}_0) = 2\text{Vol}(\mathcal{B}_{\mathcal{A}}) - 2\text{Vol}(\text{Conv}_{\text{hull}}(\mathcal{B}_{\mathcal{A}}) - \text{Conv}_{\text{hull}}(\mathcal{A}_0)) = 2\text{Vol}(\mathcal{B}_{\mathcal{A}}) - 2\text{Vol}(\mathcal{Q}).$$

Since $\mathcal{C}_{i,j}^{\mathcal{A}}$'s are disjoint in $\Delta_{\mathcal{B}_{\mathcal{A}}}$, the Dixon matrix is exact if and only if

$$2\text{Vol}(\mathcal{Q}) = 2\text{Vol}(\mathcal{Q}_{00}^{\mathcal{A}}) + 2\text{Vol}(\mathcal{Q}_{01}^{\mathcal{A}}) + 2\text{Vol}(\mathcal{Q}_{10}^{\mathcal{A}}) + 2\text{Vol}(\mathcal{Q}_{11}^{\mathcal{A}}) = |\mathcal{C}_{00}^{\mathcal{A}}| + |\mathcal{C}_{01}^{\mathcal{A}}| + |\mathcal{C}_{10}^{\mathcal{A}}| + |\mathcal{C}_{11}^{\mathcal{A}}|.$$

Proposition 5.1 (i) $|\mathcal{C}_{i,j}^{\mathcal{A}}| \leq 2\text{Vol}(\mathcal{Q}_{i,j}^{\mathcal{A}})$, and (ii) $|\mathcal{C}_{i,j}^{\mathcal{A}}| = 2\text{Vol}(\mathcal{Q}_{i,j}^{\mathcal{A}})$ if and only if each $\mathcal{C}_{i,j}^{\mathcal{A}}$ is a rectangle.

Proof: If $\mathcal{C}_{i,j}^{\mathcal{A}}$ is a rectangle, then $\mathcal{Q}_{i,j}^{\mathcal{A}}$ is a triangle, whose area is half of the area of $\mathcal{C}_{i,j}^{\mathcal{A}}$. But the number of lattice points in this rectangle belonging to $\mathcal{C}_{i,j}^{\mathcal{A}}$ equals twice the volume of $\mathcal{Q}_{i,j}^{\mathcal{A}}$.

Assume some $\mathcal{C}_{i,j}^{\mathcal{A}}$ is not a rectangle. Then, it can be broken up into vertical rectangles R_1, \dots, R_n , such that

$$\mathcal{C}_{i,j}^{\mathcal{A}} = \bigcup_{k=1}^n R_k \quad \text{and} \quad |\mathcal{C}_{i,j}^{\mathcal{A}}| = \sum_{k=1}^n |R_k|.$$

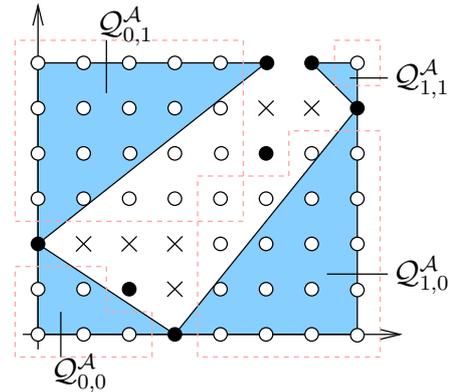


Figure 4: Newton polytope Complement

Note that $|R_i| \leq 2\text{Vol}(\text{Conv}(R_i) - \text{Conv}(\mathcal{A}))$; further $|R_i| = 2\text{Vol}(\text{Conv}(R_i) - \text{Conv}(\mathcal{A}))$ only when \mathcal{A} intersects R_i diagonally. Since \mathcal{A} cannot intersect every R_i diagonally, without making $\mathcal{C}_{i,j}^{\mathcal{A}}$ a rectangle, the number of points in $\mathcal{C}_{i,j}^{\mathcal{A}}$ is less than twice the volume of the given corner. \square .

We have the main result of the paper:

Theorem 5.2 *The size of the Dixon matrix of an unmixed generic polynomial system $\mathcal{F} = \{f_0, f_1, f_2\}$ with support \mathcal{A} is*

$$|\Delta_{\mathcal{A}}| = \mu(\mathcal{A}, \mathcal{A}) + \sum_{i,j \in \{0,1\}} (2\text{Vol}(\mathcal{Q}_{i,j}^{\mathcal{A}}) - |\mathcal{C}_{i,j}^{\mathcal{A}}|) = \mu(\mathcal{A}, \mathcal{A}) + D_e.$$

And, D_e is an upper bound on the degree of the extraneous factor in the projection operator expressed in the coefficients of f_0 , f_1 and f_2 and extracted from the Dixon matrix.

In [CK00a], a method based on partitioning the support of an unmixed polynomial system is given for estimating the degree of the extraneous factor in the projection operator extracted from the associated Dixon matrix. The above theorem generalizes that result; instead of breaking up the support into smaller supports, it gives a better insight into the existence of extraneous factors. Further, the estimate on the degree of an extraneous factor can be calculated efficiently using the above relation.

From the above theorem, there is a nice characterization of all bivariate unmixed polynomial systems for which the Dixon method computes the resultant exactly.

Theorem 5.3 *Given an unmixed generic polynomials system with support \mathcal{A} such that $\mathcal{C}_{\mathcal{A}} = \{\mathcal{C}_{00}^{\mathcal{A}}, \mathcal{C}_{01}^{\mathcal{A}}, \mathcal{C}_{10}^{\mathcal{A}}, \mathcal{C}_{11}^{\mathcal{A}}\}$, the Dixon method computes its resultant exactly if and only if each $\mathcal{C}_{i,j}^{\mathcal{A}}$ is a rectangle.*

In contrast to the results in [Chi01], Theorem 5.3 thus provides a necessary and sufficient condition on the support of an unmixed generic bivariate polynomial system for which the Dixon method computes the resultant exactly. Furthermore, Theorem 5.2 also gives an estimate of the degree of the extraneous factor in the projection operator computed by the Dixon method if an unmixed generic bivariate polynomial system does not satisfy this condition. These results are thus strict generalizations of the results in [Chi01].

Another implication of the above theorems together with Theorem 4.2 is that inclusion of terms corresponding to support-interior points in a polynomial system do not change the support of the Dixon polynomial and hence, the size of Dixon matrix and the degree of projection operator. However, inclusion of terms corresponding to points in the convex hull of the support but which are not support-interior, can contribute to the extraneous factors in the projection operator. But that is not the only source of extraneous factors in a projection operator. Even polynomial systems whose support does not have any points inside its convex hull can have extraneous factors in the projection operator computed by the Dixon method; consider example 8, for instance, in section 8 where examples are discussed.

5.1 Estimating the Degree of the Extraneous Factor

If the points in \mathcal{A} are assumed to be sorted, there is a linear algorithm to test whether or not a support is exact. We need to compute $|\mathcal{C}_{00}^{\mathcal{A}}| + |\mathcal{C}_{01}^{\mathcal{A}}| + |\mathcal{C}_{10}^{\mathcal{A}}| + |\mathcal{C}_{11}^{\mathcal{A}}|$ as well the volume of \mathcal{A} . To compute $|\mathcal{C}_{00}^{\mathcal{A}}|$, for instance, consider the following algorithm.

1. Sort points of \mathcal{A} into $\{\alpha_1, \dots, \alpha_n\}$, where each $\alpha_i = (\alpha_{i,x}, \alpha_{i,y})$, so that $\alpha_{i,x} < \alpha_{i+1,x}$ and $\alpha_{i,y} > \alpha_{i+1,y}$, the points which do not meet both conditions, are skipped; more over α_1 has the smallest possible y coordinate.

2. Then

$$|\mathcal{C}_{00}^{\mathcal{A}}| = \sum_{i=1}^{n-1} (\alpha_{i+1,x} - \alpha_{i,x}) \alpha_{i,y}.$$

Figure 5 shows how points in $\mathcal{C}_{00}^{\mathcal{A}}$ are broken up into rectangles, and their size adds up in step 2.

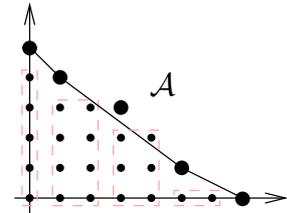


Figure 5: Computing size of $\mathcal{C}_{00}^{\mathcal{A}}$

Similar methods can be given to compute $|\mathcal{C}_{01}^A|$, $|\mathcal{C}_{10}^A|$, and $|\mathcal{C}_{11}^A|$; starting points and appropriate sorting directions are different.

Thus, for an unmixed bivariate polynomial system, it can be predicted exactly from the support, whether or not the Dixon method computes the resultant exactly, and if not, what the degree of the extraneous factor is. It should also be possible to say what coefficients of the terms in the polynomial appear in the extraneous factor as well as their degree.

6 Dixon Multiplier Matrix

As the reader would have noticed, the Dixon matrix above has, in general, complex entries; unlike in the Sylvester, Macaulay and sparse resultant formulations, where matrix entries are either zeros or coefficients of terms appearing in a polynomial system, entries in the Dixon matrix are determinants of the coefficients. For the bivariate case, entries are 3×3 determinants.

In [CK00b], we settled an open problem about the Dixon resultant formulation. We proposed a method for constructing Sylvester-type resultant matrices based on the Dixon formulation. Below, we review a generalization of that construction which has been recently developed; more details can be found in [CK02a]. We also show a relationship between these matrices and the Dixon matrices. The results in the previous section about the relationship between the support of the Dixon polynomial and the support of the polynomial system can be applied to the size of the Dixon multiplier matrices case as well.

The construction below is given first for the general case of d variables. In general, let m be any monomial. The Dixon polynomial can be rewritten as

$$m \theta(f_0, \dots, f_d) = \sum_{i=0}^d f_i \theta_i(m) \quad \text{where} \quad \theta_i(m) = \theta(f_0, \dots, f_{i-1}, m, f_{i+1}, \dots, f_d) = \overline{X} \Theta_i X_i.$$

In section 3, the Dixon polynomial was expressed through the Dixon matrix as $\theta(f_0, \dots, f_d) = \overline{X} \Theta X$. Putting both expressions for the Dixon polynomial together, we get

$$\begin{aligned} m \theta(f_0, \dots, f_d) &= m \overline{X} \Theta X = \overline{X} \Theta mX = \sum_{i=0}^d f_i \theta_i(m) = \sum_{i=0}^d f_i \overline{X} \Theta_i X_i \\ &= \overline{X} \sum_{i=0}^d \Theta_i (X_i f_i) = \overline{X} (\Theta_0 : \Theta_1 : \dots : \Theta_d) \begin{pmatrix} X_0 f_0 \\ X_1 f_1 \\ \vdots \\ X_d f_d \end{pmatrix}. \end{aligned}$$

The above matrix can also be written as:

$$\overline{X} (T \times M) X',$$

where $T = (\Theta_0 : \Theta_1 : \dots : \Theta_d)$ and $M \times X'$ is the Dixon multiplier matrix constructed from $(X_i f_i)$ for $i = \{0, \dots, d\}$.

It is proved in [CK00b] that for the special case of $m = 1$, the matrix M is a Sylvester-type resultant matrix with entries 0 and coefficients of terms in polynomials in \mathcal{F} . Further, a projection operator can be extracted as the determinant of a rank submatrix of M [KSY94].³ The matrix T is called the *transformation* matrix, and it relates the Dixon matrix to the associated Sylvester-type matrix (called the sparse Dixon matrix in [CK00b] and called the Dixon multiplier matrix in this paper). The set X_i of terms is a multiplier set for f_i .

³In [CK00b], the monomial 1 is used for the construction of the Dixon multiplier matrices, which are called **sparse Dixon** matrices in [CK00b]. The above construction is a generalization of the construction in [CK00b]. This generalization turns out to be particularly useful for constructing “good” Dixon multiplier matrices for mixed polynomial systems; the determinants of such multiplier matrices have smaller degree extraneous factors in the associated projection operators.

It is shown in [CK02a] that for an unmixed polynomial system \mathcal{F} with support \mathcal{A} , if $S_{\mathbf{x}}(m) \in \overline{\mathcal{A}}$ (recall that $\overline{\mathcal{A}}$ is the union of \mathcal{A} and all support-interior points of \mathcal{A}), then

$$S_{\mathbf{x}}(\theta) = S_{\mathbf{x}}(\theta_i(m)) \quad \text{for } i = 0, \dots, d.$$

Further, the multiplier set X_i for f_i is obtained from $\theta_i(m)$; in fact, $S_{\mathbf{x}}(X_i) = S_{\mathbf{x}}(\theta_i(m))$. For the unmixed case, $S_{\mathbf{x}}(X_i)$ is the same for each $i = 0, \dots, d$; it is denoted by \mathcal{X} . Thus, in the case of a generic unmixed polynomial system, any m from $\overline{\mathcal{A}}$ can be used for constructing a Dixon multiplier matrix. For convenience, the least degree monomial m in $\overline{\mathcal{A}}$ is picked. In section 7, where mixed polynomial systems are discussed, it is shown that the choice of m becomes crucial for generating **good** Dixon multiplier matrices leading to resultants or projection operators with extraneous factors of low degree.

It is shown in [CK00b] and [CK02a] that

Theorem 6.1 *For a given polynomial system \mathcal{F} , if its Dixon matrix is exact (in the sense, that the resultant of \mathcal{F} is the determinant of the Dixon matrix), then the associated Dixon multiplier matrix is exact as well.*

For the bivariate case, the determinant of the Dixon matrix is the same, irrespective of the variable ordering used in constructing the Dixon polynomial. It is, however, possible to construct two different Dixon multiplier matrices based on different variable orderings. The two multiplier sets are:

$$\mathcal{X}_1 = \Delta_{\mathcal{B}_A}^{x,y} / \langle \mathcal{C}_{00}^A, \mathcal{C}_{01}^A, \mathcal{C}_{10}^A, \mathcal{C}_{11}^A \rangle \quad \text{or} \quad \mathcal{X}_2 = \Delta_{\mathcal{B}_A}^{y,x} / \langle \mathcal{C}_{00}^A, \mathcal{C}_{01}^A, \mathcal{C}_{10}^A, \mathcal{C}_{11}^A \rangle,$$

where $\Delta_{\mathcal{B}_A}^{x,y} = \Delta_{\mathcal{B}_A}$, as discussed in the previous section, and $\Delta_{\mathcal{B}_A}^{y,x}$ is the set constructed in the same way as $\Delta_{\mathcal{B}_A}$ except that the role of x and y is reversed.

For the unmixed generic bivariate case, if the size of the multiplier set $|\mathcal{X}| = \mu(\mathcal{A}, \mathcal{A}) = 2\text{Vol}(\mathcal{A})$, then the Dixon matrix is exact, implying that its determinant is the resultant. In that case, the Dixon multiplier matrix is also exact.

From the above theorem and Theorem 5.3 in the previous section, we have another main result of the paper:

Theorem 6.2 *For a generic unmixed bivariate polynomial system \mathcal{F} with support \mathcal{A} and support complement $\mathcal{C}_A = \mathcal{C}_{00}^A \cup \mathcal{C}_{01}^A \cup \mathcal{C}_{10}^A \cup \mathcal{C}_{11}^A$ w.r.t. \mathcal{B}_A , the Dixon multiplier matrix M is exact if and only if each $\mathcal{C}_{i,j}^A$ is a rectangle.*

Since support interior points do not play any role in determining the support of a Dixon-polynomial, as well from the fact that in general $\mathcal{C}_{\overline{\mathcal{A}}} = \mathcal{C}_A$, we get the following corollary of the above theorem.

Corollary 6.2.1 *For a generic unmixed bivariate polynomial system with support \mathcal{A} , let $\mathcal{C}_{\overline{\mathcal{A}}} = \mathcal{C}_{00}^{\overline{\mathcal{A}}} \cup \mathcal{C}_{01}^{\overline{\mathcal{A}}} \cup \mathcal{C}_{10}^{\overline{\mathcal{A}}} \cup \mathcal{C}_{11}^{\overline{\mathcal{A}}}$ be the support complement of $\overline{\mathcal{A}}$ (which includes \mathcal{A} as well as all of its support interior points) with respect to \mathcal{B}_A . The Dixon Multiplier matrix is exact if and only if each $\mathcal{C}_{i,j}^{\overline{\mathcal{A}}}$ is rectangular.*

6.1 Zhang and Goldman's Corner Cut Supports

In [ZG00, Zha00], Zhang and Goldman proposed a method to construct Sylvester-type matrices for the bivariate case. Below, we show how their results follow from our general result above. As will be shown below, our result is stronger since it gives a necessary and sufficient condition on bivariate supports.

Zhang and Goldman [ZG00] defined a *corner-cut* support as a support obtained from a bi-degree support after removing rectangular corners.

Definition 6.1 *A support \mathcal{A} is called **corner-cut** if $\mathcal{A} = \mathcal{B}_A / \langle \mathcal{C}_{00}^A, \mathcal{C}_{01}^A, \mathcal{C}_{10}^A, \mathcal{C}_{11}^A \rangle$ and each $\mathcal{C}_{i,j}^A$ is a rectangle.*

Note that above definition requires that not only each $\mathcal{C}_{i,j}^A$ is a rectangle, but also \mathcal{A} contains all of the support interior points.

For an unmixed bivariate polynomial system with a corner-cut support \mathcal{A} , Zhang and Goldman proposed to use the following multipliers to construct the resultant matrix

$$\mathcal{X} = \Delta_{\mathcal{B}_A}^{y,x} / \langle \mathcal{C}_{00}^A, \mathcal{C}_{01}^A, \mathcal{C}_{10}^A, \mathcal{C}_{11}^A \rangle.$$

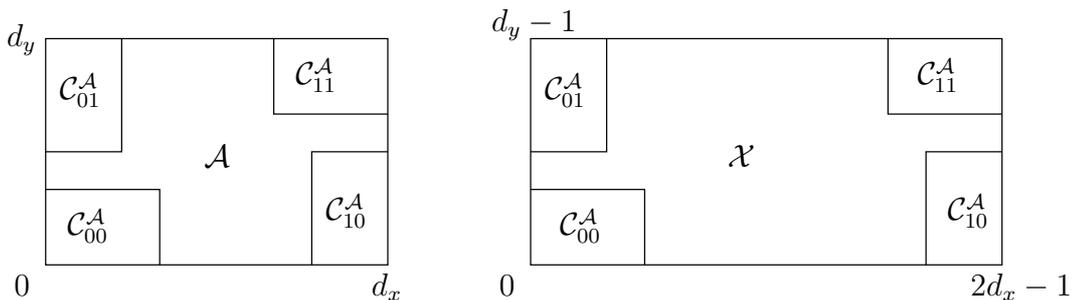


Figure 6: Corner cut support \mathcal{A} and multiplier set \mathcal{X} as in [Zha00]

In Figure 6, the support \mathcal{A} and the multiplier set \mathcal{X} used by Zhang and Goldman are shown. The Minkowski sum (whose points correspond to the columns of the resultant matrix) is shown in Figure 7. In particular,

$$\begin{aligned} |\mathcal{X}| &= |\Delta_{\mathcal{B}_A}^{y,x}| - |C_{00}^A| - |C_{01}^A| - |C_{10}^A| - |C_{11}^A| = 2 \text{Vol}(\mathcal{A}) \quad \text{and} \\ |\mathcal{A} + \mathcal{X}| &= 3|\Delta_{\mathcal{B}_A}| - 3|C_{00}^A| - 3|C_{01}^A| - 3|C_{10}^A| - 3|C_{11}^A| = 3|\mathcal{X}|. \end{aligned}$$

The matrix defined by Zhang and Goldman's construction is square, and its size is exact in the sense that each polynomial appears in the matrix as many times as the number of toric roots of the other 2 polynomials. It was shown in [ZG00] that these matrices are nonsingular in the generic case. Hence, their determinant is the resultant.

A corner-cut support \mathcal{A} satisfies the condition in Theorem 6.2, giving

Corollary 6.2.2 *Given a generic unmixed polynomial system \mathcal{F} with a corner-cut support \mathcal{A} , the determinant of the Dixon multiplier matrix constructed using multipliers from \mathcal{X} is its resultant.*

It is also possible to use the multipliers

$$\mathcal{X}' = \Delta_{\mathcal{B}_A}^{x,y} / \langle C_{00}^A, C_{01}^A, C_{10}^A, C_{11}^A \rangle,$$

giving the exact resultant.

The condition in Corollary 6.2.1 is weaker than the one required by Zhang and Goldman. Even if the support \mathcal{A} of a generic unmixed bivariate polynomial system is not corner-cut, but the support $\overline{\mathcal{A}}$ which includes \mathcal{A} and its support-interior points is corner-cut, even then the resultant can be computed exactly using the Dixon multiplier matrix construction. Furthermore, this is a necessary and sufficient condition for the determinant of the associated Dixon multiplier matrix to be the resultant.

Another immediate corollary of this result is that if \mathcal{A} is such that $\overline{\mathcal{A}}$ is not corner-cut, the determinant of the Dixon multiplier matrix constructed using multipliers from \mathcal{X} (or \mathcal{X}') is a nontrivial multiplier of its resultant (in other words, there is an extraneous factor).

The notion of a corner-cut support cannot be naturally extended to polynomial systems with more than two variables, as corner cut construction does not yield exact matrices for some simple 3 dimensional supports. For example, an unmixed polynomial system with support $\{(0, 0, 0), (1, 0, 0), (0, 1, 0), (0, 0, 2), (1, 0, 2), (0, 1, 2)\}$ can be thought of as corner-cut, as the rectangular corner is missing at points $(1, 1, 0), (1, 1, 1)$ and $(1, 1, 2)$, yet the Dixon multiplier matrix will result in extraneous factors, if care is not taken to choose appropriate variable order while constructing Dixon polynomial. This raises an interesting open question: given a corner-cut support in 3 dimensions (generalized in the natural way), does there always exist a variable order making the Dixon multiplier matrix exact?

7 Mixed Polynomial Systems

We show below that the Dixon multiplier matrix construction is especially effective for mixed polynomial systems. This construction depends upon on generating multiplier sets for each polynomial, and is determined

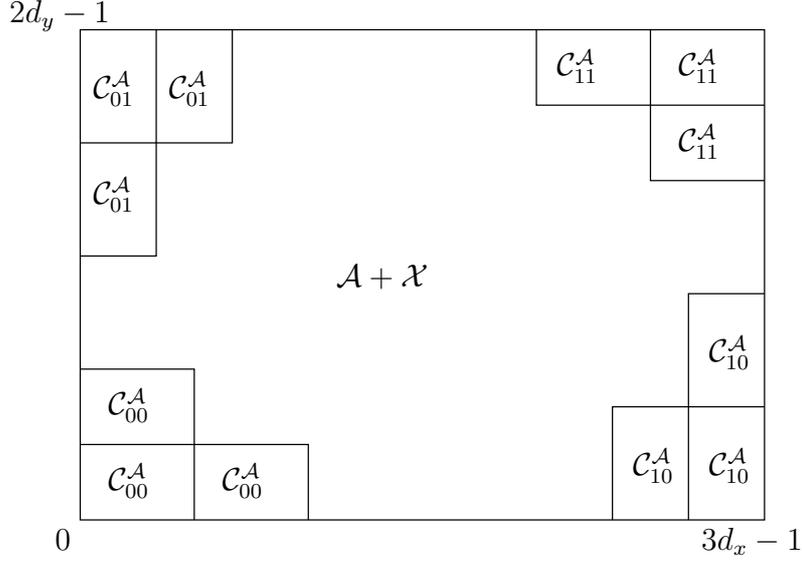


Figure 7: The Minkowski Sum $\mathcal{A} + \mathcal{X}$ as in [Zha00]

by the presence (or absence) of monomials in the support of the polynomials in the polynomial system. As should be evident from the above discussion, the multiplier sets determine the size of the Dixon multiplier matrix and hence, the degree of a projection operator.

Consider a mixed, generic polynomial bivariate system $\mathcal{F} = \{f_0, f_1, f_2\}$:

$$f_0 = \sum_{\alpha \in \mathcal{A}_0} a_{i,j} x^{\alpha_x} y^{\alpha_y}, \quad f_1 = \sum_{\beta \in \mathcal{A}_1} b_{i,j} x^{\beta_x} y^{\beta_y}, \quad f_2 = \sum_{\gamma \in \mathcal{A}_2} c_{i,j} x^{\gamma_x} y^{\gamma_y},$$

where \mathcal{A}_i is a support of f_i .

In mixed cases, the construction of a **good** Dixon multiplier matrix (in the sense that its determinant gives the projection operator that has the least degree extraneous factor) is sensitive to the choice of m in the construction which influences the support of $\theta_i(m)$ and hence, the multiplier set for each f_i . Choosing an appropriate m can be formulated as an optimization problem.

For any given support point $\alpha = (\alpha_x, \alpha_y)$, let $m = x^{\alpha_x} y^{\alpha_y}$; then

$$S_{\mathbf{x}}(\theta_0(m)) = \Delta_{\langle \{\alpha\}, \mathcal{A}_1, \mathcal{A}_2 \rangle}, S_{\mathbf{x}}(\theta_1(m)) = \Delta_{\langle \mathcal{A}_0, \{\alpha\}, \mathcal{A}_2 \rangle}, S_{\mathbf{x}}(\theta_2(m)) = \Delta_{\langle \mathcal{A}_0, \mathcal{A}_1, \{\alpha\} \rangle}.$$

Let $\Phi_0(\alpha) = |\Delta_{\langle \{\alpha\}, \mathcal{A}_1, \mathcal{A}_2 \rangle}|$, $\Phi_1(\alpha) = |\Delta_{\langle \mathcal{A}_0, \{\alpha\}, \mathcal{A}_2 \rangle}|$ and $\Phi_2(\alpha) = |\Delta_{\langle \mathcal{A}_0, \mathcal{A}_1, \{\alpha\} \rangle}|$.

Since $\Phi_i(\alpha)$ represents the number of rows corresponding to polynomial f_i in the Dixon multiplier matrix, the goal is to find α such that

$$\Phi(\alpha) = \Phi_0(\alpha) + \Phi_1(\alpha) + \Phi_2(\alpha)$$

is minimized; that is, the size of the entire Dixon matrix is minimized, so that smaller extraneous factor is produced. One can choose to minimize particular $\Phi_i(\alpha)$ in the hope of having coefficients of f_i appearing with smallest degree in the projection operator.

An α that minimizes $\Phi(\alpha)$ must belong to $\mathcal{B}_{\mathcal{A}}$, as the size $\Delta_{\mathcal{A}}$ will only increase for an α outside $\mathcal{B}_{\mathcal{A}}$.

7.1 Searching the Appropriate Monomial for Constructing Multiplier Matrix

A brute force method for finding an exponent vector that minimizes $\Phi(\alpha)$ is to compute its value for each point in $\mathcal{B}_{\mathcal{A}}$. One disadvantage of this method is that it depends upon the maximum and minimum values of exponents in the supports, which is (almost) the same as the the degree of variables of the polynomials in

the system.⁴ Even though this search is not likely to significantly affect the overall complexity of the resultant algorithm, it will be nice to have a method that depends on the size of the supports; in that case, polynomial systems with small supports can be easily handled.

We show below that the bounding box \mathcal{B}_A of supports can be partitioned into rectangular regions, called *cells*, such that each $\Phi_i(\alpha)$ is linear in the coordinates of $\alpha = (\alpha_x, \alpha_y)$ over a cell. The equation for each $\Phi_i(\alpha)$ changes from one cell to another.

Let $P = \mathcal{A}_0 \cup \mathcal{A}_1 \cup \mathcal{A}_2$ and $P = \langle p_{i_1}, \dots, p_{i_n} \rangle$ be the list of points sorted on the x coordinates, and similarly, $P' = \langle p_{j_1}, \dots, p_{j_n} \rangle$ be the list of points sorted on the y coordinates.

Let $\mathcal{B}_A = \mathcal{B}_u^v$. Split the values of the x coordinates into $2n + 1$ intervals I_1, \dots, I_{2n+1} ,

$$(u_x, p_{i_1, x}), [p_{i_1, x}, p_{i_1, x}], (p_{i_1, x}, p_{i_2, x}), \dots, (p_{i_{n-1}, x}, p_{i_n, x}), [p_{i_n, x}, p_{i_n, x}], (p_{i_n, x}, v_x),$$

and similarly, split the values of the y coordinates into $2n + 1$ intervals J_1, \dots, J_{2n+1} ,

$$(u_y, p_{j_1, y}), [p_{j_1, y}, p_{j_1, y}], (p_{j_1, y}, p_{j_2, y}), \dots, (p_{j_{n-1}, y}, p_{j_n, y}), [p_{j_n, y}, p_{j_n, y}], (p_{j_n, y}, v_y),$$

where $p_{k,x}, p_{k,y}$ are, respectively the x and y coordinates of p_k , (a, b) is an open interval containing integral points strictly greater than a and strictly less than b , and $[a, b]$ is a closed interval containing integral points greater than or equal to a and less than or equal to b . This splits \mathcal{B}_A into rectangular regions $R_{k,l}(P) = I_k \times J_l$ for $k, l \in \{1, \dots, 2n + 1\}$. Note that every point of \mathcal{B}_u^v belongs to only one region. Let $R_{k,l}(P)$ stand for the $(k, l)^{th}$ region of \mathcal{B}_u^v , based on P .

Theorem 7.1 $\Phi(\alpha)$ is linear in cell $R_{k,l}(P)$ for any $k, l \in \{1, \dots, 2n + 1\}$.

Proof: (Sketch) It suffices to show that each $\Phi_i(\alpha)$ is linear in $R_{k,l}(P)$. Below, we sketch the proof of linearity of Φ_0 ; proofs for Φ_1 and Φ_2 are similar.

From the definition of $R_{k,l}(P)$, it should be evident that the relative position of any $\alpha \in R_{k,l}(P)$ to the points of \mathcal{A}_1 and \mathcal{A}_2 does not change. One way to establish linearity is to subdivide all simplexes $\langle \alpha, \beta, \gamma \rangle$ where $\beta \in \mathcal{A}_1$ and $\gamma \in \mathcal{A}_2$ as in Proposition 4.2, which also can be done on the y coordinates. Since α is restricted to $R_{k,l}(P)$, most of the simplexes will not even change when α varies within the cell. It can be shown that rectangular regions generated by these simplexes will increase or decrease by the same amount as the coordinates of α are changed \square .

Given that the minimal value of a linear equation over a polyhedral is always at corners, it suffices to evaluate $\Phi(\alpha)$ at the corner of each cell.

The complexity of this algorithm is governed by the number of cells and the cost of evaluating $\Phi(\alpha)$. Number of cells is bounded by $(2n + 1)^2$, where n is the size of \mathcal{A} , the union of all supports. The evaluation cost of $\Phi(\alpha)$ is the sum of the evaluation cost of each $\Phi_i(\alpha)$ for $i = \{0, 1, 2\}$. The evaluation cost of each $\Phi_0(\alpha)$ is $O(n_1 n_2)$ where $n_1 = |\mathcal{A}_1|, n_2 = |\mathcal{A}_2|$. The overall complexity for finding an α that minimizes $\Phi(\alpha)$ is

$$O((2n + 1)^2(n_1 n_2 + n_0 n_2 + n_0 n_1)) = O(n^4).$$

As an example, consider the polynomial system shown in example 4 of the next section. Its support is:

$$\begin{aligned} \mathcal{A}_0 &= \{(1, 0), (0, 1)\}, \\ \mathcal{A}_1 &= \{(1, 1), (1, 2)\}, \quad \text{and} \\ \mathcal{A}_2 &= \{(2, 1), (0, 1)\}. \end{aligned}$$

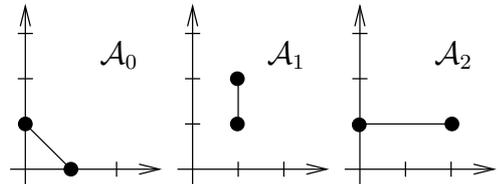


Figure 8: Support of example 4

The set \mathcal{B}_A has all the points with the maximum coordinate value of 2. Table 1 below shows all values $\Phi(\alpha)$ for $\alpha \in \mathcal{B}_A$, as well for α outside \mathcal{B}_A . Since this problem is easy for a Gröbner basis algorithm, its toric resultant can be obtained using the Gröbner basis computation:

$$gb = a_{10}^2 c_{01} b_{12}^2 + c_{21} a_{01}^2 b_{11}^2.$$

⁴It is the same if there is a constant term in at least one polynomial.

It is thus possible to compare the extraneous factors when different α 's are chosen.

No	α	$\Phi_0(\alpha)$	$\Phi_1(\alpha)$	$\Phi_2(\alpha)$	$\Phi(\alpha)$	nCols	Rank	Extra Factor
1.	(0, 0)	4	3	3	10	10	9	$b_{11}a_{10}^2c_{21}$
2.	(1, 0)	4	2	2	8	8	8	$a_{10}^2c_{01}$
3.	(0, 1)	2	2	2	6	7	6	$-c_{21}$
4.	(2, 0)	4	3	4	11	11	10	$-a_{01}a_{10}c_{01}c_{21}^2$
5.	(1, 1)	2	2	1	5	5	5	1
6.	(0, 2)	3	4	2	9	8	8	$a_{01}b_{12}c_{21}$
7.	(3, 0)	5	4	6	15	14	12	$-a_{01}^3c_{21}^3c_{01}$
8.	(2, 1)	2	2	3	7	8	7	$-c_{21}^2$
9.	(1, 2)	2	4	1	7	7	7	$-b_{12}^2$
10.	(0, 3)	5	6	3	14	11	11	$a_{01}^2a_{10}b_{11}b_{12}c_{21}$
11.	(4, 0)	6	5	8	19	17	15	$-a_{01}^4b_{11}c_{21}^4c_{01}$
12.	(3, 1)	3	3	5	11	12	10	$-c_{21}^3b_{11}a_{01}$
13.	(2, 2)	3	4	3	10	10	10	$-b_{12}^2c_{21}^2a_{10}$
14.	(1, 3)	4	6	2	12	10	10	$-b_{11}b_{12}^2a_{01}a_{10}$
15.	(0, 4)	7	8	4	19	14	14	$a_{01}^4b_{11}b_{12}c_{21}^2a_{10}$

Table 1: Choice of α for the Dixon multiplier of example 4.

The size of the exact resultant matrix is 5 as $\mu(\mathcal{A}_1, \mathcal{A}_2) = 2, \mu(\mathcal{A}_0, \mathcal{A}_2) = 2$ and $\mu(\mathcal{A}_0, \mathcal{A}_1) = 1$; hence, optimal $\Phi_0(\alpha) = 2, \Phi_1(\alpha) = 2$ and $\Phi_2(\alpha) = 1$. From the table $\alpha = (1, 1)$ satisfies this condition.

Note that in the extraneous factor column, the degree of the coefficients from f_i is $\Phi_i(\alpha) - \mu_i(\mathcal{A})$ in general, for $i = 0, 1, 2$, where $\mu_0(\mathcal{A}) = \mu(\mathcal{A}_1, \mathcal{A}_2)$, $\mu_1(\mathcal{A}) = \mu(\mathcal{A}_0, \mathcal{A}_2)$ and $\mu_2(\mathcal{A}) = \mu(\mathcal{A}_0, \mathcal{A}_1)$. In case the matrix is singular, the degree of the extraneous factor can be smaller, depending upon the maximal minor selected from the resultant matrix.

8 Examples

In this section, we discuss a family of bivariate systems discussed in the literature [DE01a], and the performance of different algorithms in generating resultant matrices. Details of the examples are given after the table. The column *deg R* gives the degree of the resultant. Each method is identified by the paper in which it appeared. The last two columns are for the two methods based on the Dixon formulation discussed in this paper. The column labelled $|\Delta_{\mathcal{A}}|$ gives the size of the Dixon matrix. The reader should recall that the entries in the Dixon matrix are 3×3 determinants expressed in the coefficients of the terms in the polynomials. For other methods, matrix entries are mostly zeros or coefficients of terms in the polynomials. The last column in the table is the size of the Dixon multiplier matrix. For unmixed cases, the size of the Dixon multiplier matrix is 3 times the size of the Dixon matrix.

In the last column, if there is * next to the entry, this indicates that the algorithm in section 7 for generating the Dixon multiplier matrix by choosing an appropriate monomial was used instead of treating the problem as though it was unmixed by considering the union of the supports of the polynomials.

The degree of projection operators (hence extraneous factors) cannot be determined from the matrix sizes in the case of [CDS98], [DE01a] as well as the Dixon matrix (the $\Delta_{\mathcal{A}}$ column) as some of the matrix entries are different from coefficients of terms in polynomials. For the Dixon matrix, the degree of projection operator is $3|\Delta_{\mathcal{A}}|$; for the methods in [CDS98] and [DE01a], the degree of the projection operator is 2 more than the matrix size.

From the table, it is clear that the Dixon multiplier matrix method produces smaller extraneous factors; in almost all examples, it computes the resultant exactly. The method also turns out to be computationally less expensive for extracting a projection operator.

Ex.	deg R	[CE00]	[CDS98]	[ZG00]	[DE01a]	$ \Delta_{\mathcal{A}} $	$\Phi(\alpha)$
1	$3n^2$	$\frac{3}{2}n(3n-1)$	$6n^2 - 9n + 4$	$\frac{21}{4}n^2$	$\frac{9}{2}n(n-1)+1$	$\frac{1}{2}n(3n-1)$	$\frac{3}{2}n(3n-1)$
2	$6n_1n_2$	$9n_1n_2$	$12n_1n_2 - 6(n_1+n_2)+4$	$6n_1n_2$	$9n_1n_2 - 3(n_1+n_2)+1$	$2n_1n_2$	$6n_1n_2$
3	12	15	10	12	10	4	12
4	5	5	4	6	4	2	5*
5	7	12	7	15	7	4	7*
6	18	26	22	18	19	6	18
7	24	35	34	30	22	8	24*
8	57	92	97	63	64	20	60

Table 2: Comparison of resultant matrices

1. **Homogeneous** (unmixed) polynomial system of degree n :

$$f_0(x, y) = \sum_{i+j \leq n} a_{ij}x^i y^j, \quad f_1(x, y) = \sum_{i+j \leq n} b_{ij}x^i y^j, \quad f_2(x, y) = \sum_{i+j \leq n} c_{ij}x^i y^j.$$

The mixed volume of any two polynomials is n^2 , the Bezout bound. The degree of the resultant is $3n^2$.

2. **Bi-homogeneous** (unmixed, corner cut) polynomial system of degree n_1, n_2 :

$$f_0(x, y) = \sum_{i=0}^{n_1} \sum_{j=0}^{n_2} a_{ij}x^i y^j, \quad f_1(x, y) = \sum_{i=0}^{n_1} \sum_{j=0}^{n_2} b_{ij}x^i y^j, \quad f_2(x, y) = \sum_{i=0}^{n_1} \sum_{j=0}^{n_2} c_{ij}x^i y^j.$$

The mixed volume of any two polynomials is $2n_1n_2$. The degree of the resultant is $6n_1n_2$.

3. **Examples from [CDS98]** (unmixed, corner cut):

$$\begin{aligned} f_0(x, y) &= a_{00} + a_{01}y + a_{10}x + a_{11}xy + a_{12}xy^2 + a_{13}xy^3, \\ f_1(x, y) &= b_{00} + b_{01}y + b_{10}x + b_{11}xy + b_{12}xy^2 + b_{13}xy^3, \\ f_2(x, y) &= c_{00} + c_{01}y + c_{10}x + c_{11}xy + c_{12}xy^2 + c_{13}xy^3. \end{aligned}$$

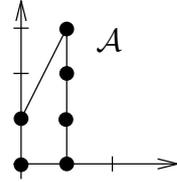


Figure 9: Support of example 3

This problem is given as an example in [CDS98] of the Chow form of a Hilzebruch surface. It is an unmixed problem, where any two polynomials have the mixed volume of 4. Notice that this problem has a corner-cut support.

4. **Example from [ZG00]** (mixed) This example appeared earlier in section 7. It served as an example in [DE01a] and [ZG00].

$$f_0(s, t) = 2s + t, \quad f_1(s, t) = st + st^2, \quad f_2 = s^2t + 2t.$$

The mixed volume is $[2, 2, 1]$ (i.e., $\mu(\mathcal{A}_1, \mathcal{A}_2) = 2$, $\mu(\mathcal{A}_0, \mathcal{A}_2) = 2$ and $\mu(\mathcal{A}_0, \mathcal{A}_1) = 1$). The degree of the resultant is 5.

5. **Example from [Man92]** (Mixed):

$$\begin{aligned} f_0(x, y) &= a_{10}x + a_{20}x^2 + a_{01}y, \\ f_1(x, y) &= b_{10}x + b_{02}y^2 + b_{01}y, \\ f_2(x, y) &= c_{10}x + c_{11}xy + c_{01}y. \end{aligned}$$

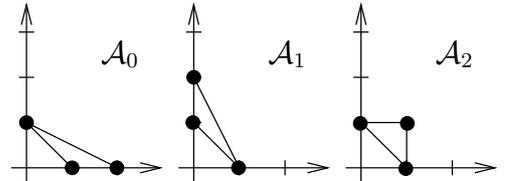


Figure 10: Support of example 5

This problem is about surface parameterization [Man92]. Its BKK bound is $[2, 2, 3] = 7$.

6. **Example from [DE01b]** (unmixed, corner cut):

$$\begin{aligned} f_0 &= a_{00} + a_{10}x + a_{01}y + a_{12}xy^2 + a_{21}x^2y + a_{22}x^2y^2, \\ f_1 &= b_{00} + b_{10}x + b_{01}y + b_{12}xy^2 + b_{21}x^2y + b_{22}x^2y^2, \\ f_2 &= c_{00} + c_{10}x + c_{01}y + c_{12}xy^2 + c_{21}x^2y + c_{22}x^2y^2. \end{aligned}$$

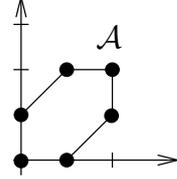


Figure 11: Support of example 6

The mixed volume of any two polynomials is 6; therefore, the degree of the resultant is 18. Moreover, the problem is unmixed and corner-cut; therefore, the Dixon method, the Dixon multiplier matrix method and [ZG00] have exact matrices for this problem.

It is included in [DE01b]; it is interesting because the hybrid method proposed in [DE01b] does not produce an exact resultant matrix.

7. **Example from [GS01]** (mixed):

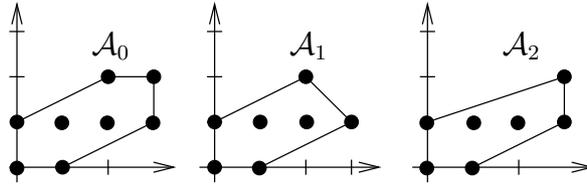


Figure 12: Support of example 7

$$\begin{aligned} f_0 &= a_{00} + a_{10}x + a_{01}y + a_{11}xy + a_{21}x^2y + a_{22}x^2y^2 + a_{31}x^3y + a_{32}x^3y^2, \\ f_1 &= b_{00} + b_{10}x + b_{01}y + b_{11}xy + b_{21}x^2y + b_{22}x^2y^2 + b_{31}x^3y, \\ f_2 &= c_{00} + c_{10}x + c_{01}y + c_{11}xy + c_{21}x^2y + c_{31}x^3y + c_{32}x^3y^2. \end{aligned}$$

This polynomial system is defined in [GS01] to study the self intersections of a parameterized surface. Interestingly, this problem has the BKK bound of $[8, 8, 8] = 24$, which is the same as though this system was unmixed whose support equals the support of the first polynomial which is also the union of the supports of the other two polynomials.

8. **Example from [DE01a]** (unmixed)

$$\begin{aligned} f_0(x, y) &= a_{10}x + a_{21}x^2y + a_{03}y^3 + a_{15}xy^5 + a_{25}x^2y^5 + a_{33}x^3y^3 + a_{34}x^3y^4, \\ f_1(x, y) &= b_{10}x + b_{21}x^2y + b_{03}y^3 + b_{15}xy^5 + b_{25}x^2y^5 + b_{33}x^3y^3 + b_{34}x^3y^4, \\ f_2(x, y) &= c_{10}x + c_{21}x^2y + c_{03}y^3 + c_{15}xy^5 + c_{25}x^2y^5 + c_{33}x^3y^3 + c_{34}x^3y^4. \end{aligned}$$

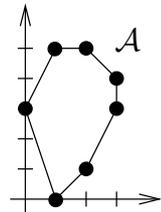


Figure 13: Support of example 8

This example appeared in [DE01a] as a demonstration for the hybrid method proposed in that paper for bivariate systems. The mixed volume of any two polynomials is 19; hence, the resultant degree is 57.

9 Generalization to Trivariate and other Multivariate Polynomial Systems

We have identified necessary and sufficient conditions on the support of an unmixed bivariate polynomial system such that the methods based on the Dixon resultant formulation can compute its resultant exactly. When this cannot be done, the degree of the projection operator can be predicted, from which the degree of the extraneous factor appearing in it can be computed. These results are thus strict generalizations of the related results reported in [Chi01, ZG00, Zha00]. These results still do not lead to precisely identifying the extraneous factor as is known in the case of eliminating a single variable. We plan to investigate this issue next.

For non-generic polynomial systems for which the number of toric roots is still the BKK bound, the degree of extraneous factor in a projection operator cannot be estimated. It appears that the discrepancy between the BKK bound and the size of Dixon matrix is due to the difference in the volume of the Newton polytope and the size of the corresponding support. Experimental evidence suggests that the coefficients of terms in polynomials also play a role in determining the support of the Dixon polynomial; this is also reflected in the formula for the Dixon polynomial based on the Cauchy-Binet formula. We are interested in analyzing whether the genericity requirements for obtaining the BKK bound is sufficient to preclude any role the coefficients of terms in a polynomial system play in determining the support of the Dixon polynomial and hence, the size of Dixon matrix.

The focus of this paper has been on bivariate systems. We are interested in generalizing these results – particularly the notion of a support-interior point in an arbitrary dimension. As illustrated above, it can be shown that the determinant of the Dixon multiplier matrix is not exactly the resultant for a tri-variate generic unmixed system even if its support is corner-cut.⁵ However, a necessary and sufficient condition on supports based on exclusion of support-interior points seems plausible.

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⁵The notion of a corner-cut support is assumed to be generalized to arbitrary dimensions in an obvious way.

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