

CS 591.03

Introduction to Data Mining

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LECTURE 9: GRAPH MINING

Graph Similarity

Edit distance/graph isomorphism:

- Tree Edit Distance

Feature extraction

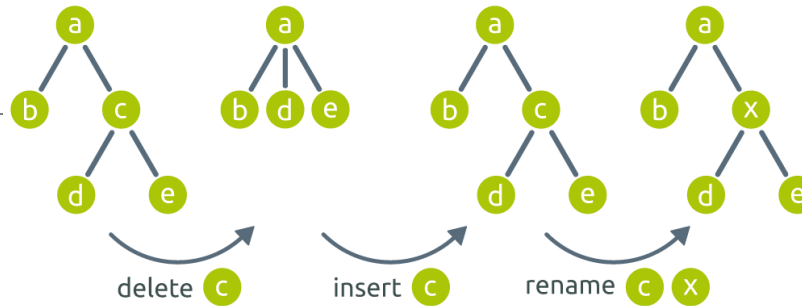
- IN/out degree
- Diameter

Iterative methods

- SimRank

Edit Distance

Three Operations



Tai's algorithm runs in $O(m^3n^3)$ time and space for trees with m and n nodes respectively.

left

$$\triangleleft, \triangleleft, \triangleleft, \triangleleft, \triangleleft = \min \left\{ \begin{array}{l} \triangleleft, \triangleleft, \triangleleft, \triangleleft + \text{del } \circ \\ \triangleleft, \triangleleft, \triangleleft, \triangleleft + \text{ins } \circ \\ \triangleleft, \triangleleft, \triangleleft + \triangleleft, \triangleleft + \text{ren } \circ \circ \end{array} \right.$$

right

$$\triangleleft, \triangleleft, \triangleleft, \triangleleft, \triangleleft = \min \left\{ \begin{array}{l} \triangleleft, \triangleleft, \triangleleft, \triangleleft + \text{del } \bullet \\ \triangleleft, \triangleleft, \triangleleft, \triangleleft + \text{ins } \bullet \\ \triangleleft, \triangleleft, \triangleleft + \triangleleft, \triangleleft + \text{ren } \bullet \bullet \end{array} \right.$$

Diameter

Largest Shortest path in the graph.

$\text{shortestPath}(i, j, 0) = w(i, j)$

$\text{shortestPath}(i, j, k+1) = \min(\text{shortestPath}(i, j, k), \text{shortestPath}(i, k+1, k) + \text{shortestPath}(k+1, j, k))$

1 let dist be a $|V| \times |V|$ array of minimum distances initialized to ∞ (infinity)

2 for each vertex v

3 $\text{dist}[v][v] \leftarrow 0$

4 for each edge (u, v)

5 $\text{dist}[u][v] \leftarrow w(u, v)$ // the weight of the edge (u, v)

6 for k from 1 to $|V|$

7 for i from 1 to $|V|$

8 for j from 1 to $|V|$

9 if $\text{dist}[i][j] > \text{dist}[i][k] + \text{dist}[k][j]$

10 $\text{dist}[i][j] \leftarrow \text{dist}[i][k] + \text{dist}[k][j]$

11 end if

Simrank

$$s(a, b) = \frac{C}{|I(a)||I(b)|} \sum_{i=1}^{|I(a)|} \sum_{j=1}^{|I(b)|} s(I_i(a), I_j(b))$$

For a node v in a graph, we denote by $I(v)$ and $O(v)$ the set of in-neighbors and out-neighbors of v , respectively.

1. A solution $s(*, *) \in [0, 1]$ to the n^2 SimRank equations always exists and is unique.
2. Symmetric
3. Reflexive

Ranking Nodes

Page Rank

$$PR(A) = (1-d) + d (PR(T1)/C(T1) + \dots + PR(Tn)/C(Tn))$$

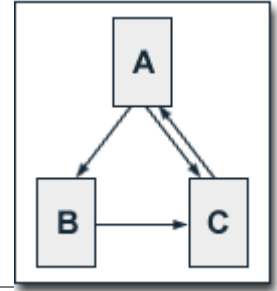
PR(A) is the PageRank of page A,

PR(Ti) is the PageRank of pages Ti which link to page A,

C(Ti) is the number of outbound links on page Ti and

d is a damping factor which can be set between 0 and 1.

Example



$$PR(A) = 0.5 + 0.5 PR(C)$$

$$PR(B) = 0.5 + 0.5 (PR(A) / 2)$$

$$PR(C) = 0.5 + 0.5 (PR(A) / 2 + PR(B))$$

These equations can easily be solved. We get the following PageRank values for the single pages:

$$PR(A) = 14/13 = 1.07692308$$

$$PR(B) = 10/13 = 0.76923077$$

$$PR(C) = 15/13 = 1.15384615$$

HITS: Hyperlink-Induced Topic Search

Iterate(G, k)

G : a collection of n linked pages

k : a natural number

Let z denote the vector $(1, 1, 1, \dots, 1) \in \mathbf{R}^n$.

Set $x_0 := z$.

Set $y_0 := z$.

For $i = 1, 2, \dots, k$

 Apply the \mathcal{I} operation to (x_{i-1}, y_{i-1}) , obtaining new x -weights x'_i .

 Apply the \mathcal{O} operation to (x'_i, y_{i-1}) , obtaining new y -weights y'_i .

 Normalize x'_i , obtaining x_i .

 Normalize y'_i , obtaining y_i .

End

Return (x_k, y_k) .