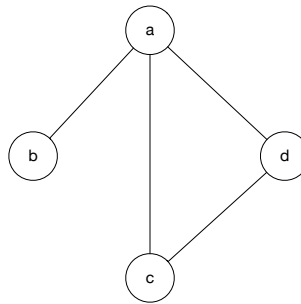


# CS 261, HW9

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*Due: May 2nd*

1. A *graph* consists of a set of vertices,  $V$ , and a set of edges  $E$ , where each edge in  $E$  is a tuple of 2 vertices in  $V$ . For example, we could have  $V = \{a, b, c, d\}$  and  $E = \{(a, b), (c, d), (a, d), (a, c)\}$ . The *degree* of a vertex  $v \in V$  is the number of edges containing  $v$ . In the example graph, the degree of  $a$  is 3, and the maximum degree of the graph is 3.



A *coloring* of a graph is an assignment of colors to all vertices in  $V$  such that for all  $(u, v) \in E$ , vertex  $u$  has a different color than vertex  $v$ . A coloring of the example graph with 3 colors is: vertex  $a$  gets color 1, vertex  $b$  gets color 2, vertex  $c$  gets color 2, and vertex  $d$  gets color 3. Prove the following.

- Any graph with maximum degree of 3 can be colored with 4 colors. Prove this by induction over  $n = |V|$ . For your IH, assume that any graph with maximum degree of 3 and **less than**  $n$  vertices can be colored with 4 colors.
2. We say that a graph  $G = (V, E)$  is *connected* if for every pair of vertices  $u, v \in V$ , there exists a path from  $u$  to  $v$  using only edges in  $G$ . To illustrate, the example graph in Problem 1 is connected, but becomes

disconnected if we remove edge  $(a, b)$ . For a graph  $G = (V, E)$  and vertex  $v \in V$ , we define  $G - v$  to be the new graph  $(V - v, E')$  where  $E'$  is  $E$  minus all edges containing  $v$ . Prove the following:

- Every connected graph  $G = (V, E)$  with  $|V| \geq 2$  has at least two vertices  $x_1$  and  $x_2$ , such that  $G - x_i$  is connected for  $i = 1, 2$ . Prove this by induction on  $n = |V|$ . For your IH, assume that any graph with  $2 < |V| < n$  has at least two vertices  $x_1$  and  $x_2$ , such that  $G - x_i$  is connected for  $i = 1, 2$ .
3. Prove by induction that  $\sum_{i=0}^n r^i = \frac{r^{n+1}-1}{r-1}$  for every  $n > 0$ , and all  $r \neq 1$ . For your IH, assume that  $\sum_{i=0}^{n-1} r^i = \frac{r^n-1}{r-1}$ .
  4. Consider the recurrence  $f(n) = 3f(n/2) + \sqrt{n}$ , where  $f(c) = \theta(1)$ , for constants  $c$ . Use the Master method to solve this recurrence.
  5. Consider the following function:

```
int f (int n){
    if (n==0) return 2;
    else if (n==1) return 5;
    else{
        int val = 2*f (n-1);
        val = val - f (n-2);
        return val;
    }
}
```

- (a) Write a recurrence relation for the *value* returned by  $f$ . Solve the recurrence exactly using annihilators. (Don't forget to check it)
- (b) Write a recurrence relation for the *running time* of  $f$ . Get a tight upperbound (i.e. big-O) on the solution to this recurrence using annihilators.