1. Consider the recurrence $T(n) = 3T(n/2) + n$

   (a) Use the recursion tree method to get a tight upper bound (i.e. big-O) on the solution to this recurrence

   (b) Now use annihilators (and a transformation) to get a tight upper bound on the solution to this recurrence. Show your work. (Note that your two bounds should match)

2. Consider the recurrence $T(n) = 2T(n/2) + \log^2 n$

   (a) Use the Master method to get a general solution to this recurrence.

   (b) Now use annihilators (and a transformation) to get a tight upper bound on the solution to this recurrence. Show your work. (Note that your two bounds should match)

3. Consider the following function:

   ```c
   int f (int n){
      if (n==0) return 3;
      else if (n==1) return 5;
      else{
         int val = 2*f (n-1);
         val = val - f (n-2);
         return val;
      }
   }
   ```

   ```c
   ```
(a) Write a recurrence relation for the value returned by \( f \). Solve the recurrence exactly. (Don’t forget to check it)

(b) Write a recurrence relation for the running time of \( f \). Get a tight upperbound (i.e. big-O) on the solution to this recurrence.

4. Give the table for the optimal string alignment of the strings “abb-baa” and “ababab”. Make sure you include arrows in your table the same way we did in the table given in lecture, these arrows will help you reconstruct the solution. List all optimal alignments of these two strings.

5. Find the optimal parenthesization for a matrix-chain product whose sequence of dimensions is: \( (2, 3, 2, 2, 1) \). (Don’t forget to include the table used to compute your result)

6. Exercise 15.2-3

7. Exercise 15.2-5 (hint: use proof by induction)

8. Exercise 15.4-1

9. Exercise 15.4-5 (note: monotonically increasing means non-decreasing, e.g. 1, 2, 2, 4, 5, 5, 7)