Quicksort

- Based on divide and conquer strategy
- Worst case is $\Theta(n^2)$
- Expected running time is $\Theta(n \log n)$
- An In-place sorting algorithm
- Almost always the fastest sorting algorithm
Quicksort

- **Divide:** Pick some element \( A[q] \) of the array \( A \) and partition \( A \) into two arrays \( A_1 \) and \( A_2 \) such that every element in \( A_1 \) is \( \leq A[q] \), and every element in \( A_2 \) is \( > A[p] \).
- **Conquer:** Recursively sort \( A_1 \) and \( A_2 \).
- **Combine:** \( A_1 \) concatenated with \( A[q] \) concatenated with \( A_2 \) is now the sorted version of \( A \).
The Algorithm

//PRE: A is the array to be sorted, p>=1;
//    r is <= the size of A
//POST: A[p..r] is in sorted order
Quicksort (A,p,r){
    if (p<r){
        q = Partition (A,p,r);
        Quicksort (A,p,q-1);
        Quicksort (A,q+1,r);
    }
}
Partition

//PRE: A[p..r] is the array to be partitioned, p>=1 and r <= size of A, A[r] is the pivot element
//POST: Let A’ be the array A after the function is run. Then
// A’[p..r] contains the same elements as A[p..r]. Further,
// all elements in A’[p..res-1] are <= A[r], A’[res] = A[r],
// and all elements in A’[res+1..r] are > A[r]

Partition (A,p,r){
    x = A[r];
    i = p-1;
    for (j=p;j<=r-1;j++){
        if (A[j]<=x){
            i++;
            exchange A[i] and A[j];
        }
    }
    exchange A[i+1] and A[r];
    return i+1;
}"
Analysis

• The function Partition takes $O(n)$ time. Why?
Example QuickSort

- QuickSort the array [2, 6, 9, 1, 5, 3, 8, 7, 4]
Randomized Quick-Sort

- We’d like to ensure that we get reasonably good splits reasonably quickly
- Q: How do we ensure that we “usually” get good splits? How can we ensure this even for worst case inputs?
- A: We use randomization.
R-Partition

//PRE: A[p..r] is the array to be partitioned, p>=1 and r <= size of A
//POST: Let A’ be the array A after the function is run. Then
//A’[p..r] contains the same elements as A[p..r]. Further,
//all elements in A’[p..res-1] are <= A[i], A’[res] = A[i],
//and all elements in A’[res+1..r] are > A[i], where i is
//a random number between $p$ and $r$.

R-Partition (A,p,r){
    i = Random(p,r);
    exchange A[r] and A[i];
    return Partition(A,p,r);
}
Randomized Quicksort

//PRE: A is the array to be sorted, p>=1, and r is <= the size of A
//POST: A[p..r] is in sorted order
R-Quicksort (A,p,r){
    if (p<r){
        q = R-Partition (A,p,r);
        R-Quicksort (A,p,q-1);
        R-Quicksort (A,q+1,r);
    }
}
Analysis

- R-Quicksort is a *randomized* algorithm
- The run time is a *random variable*
- We’d like to analyze the *expected* run time of R-Quicksort
- To do this, we first need to learn some basic probability theory.
Probability Definitions

(from Appendix C.3)

- A random variable is a variable that takes on one of several values, each with some probability. (Example: if $X$ is the outcome of the roll of a die, $X$ is a random variable)
- The expected value of a random variable, $X$ is defined as:

$$E(X) = \sum_x x \Pr(X = x)$$

(Example if $X$ is the outcome of the roll of a three sided die,

$$E(X) = 1(1/3) + 2(1/3) + 3(1/3)$$

$$= 2$$

11
Probability Definitions

• Two events $A$ and $B$ are mutually exclusive if $A \cap B$ is the empty set (Example: $A$ is the event that the outcome of a die is 1 and $B$ is the event that the outcome of a die is 2).

• Two random variables $X$ and $Y$ are independent if for all $x$ and $y$, $P(X = x \text{ and } Y = y) = P(X = x)P(Y = y)$ (Example: let $X$ be the outcome of the first roll of a die, and $Y$ be the outcome of the second roll of the die. Then $X$ and $Y$ are independent.)
**Indicator Random Variables**

- An *Indicator Random Variable* for event $A$, is defined as:

$$I(A) = \begin{cases} 
1 & \text{if event } A \text{ occurs} \\
0 & \text{otherwise} 
\end{cases}$$

- Example: Let $A$ be the event that the roll of a die equals 2. Then $I(A)$ is 1 if the die roll is 2 and 0 otherwise.
Linearity of Expectation

- Let $X$ and $Y$ be two random variables
- Then $E(X + Y) = E(X) + E(Y)$
- (Holds even if $X$ and $Y$ are not independent.)

- More generally, let $X_1, X_2, \ldots, X_n$ be $n$ random variables
- Then

$$E \left( \sum_{i=1}^{n} X_i \right) = \sum_{i=1}^{n} E(X_i)$$
Example

- For $1 \leq i \leq n$, let $X_i$ be the outcome of the $i$-th roll of a three-sided die.
- Then

$$E\left(\sum_{i=1}^{n} X_i\right) = \sum_{i=1}^{n} E(X_i) = 2n$$
Example

• Indicator Random Variables and Linearity of Expectation used together are a very powerful tool
• The *Birthday Paradox* illustrates this point
• To analyze the run time of Quicksort, we will also use indicator r.v.’s and linearity of expectation (analysis will be similar to “birthday paradox” problem)
Birthday Paradox

- Assume there are $m$ people in a room, and $n$ days in a year
- Assume that each of these $m$ people is born on a day chosen independently and uniformly at random from the $n$ days
- Q: What is the expected number of pairs of individuals that have the same birthday?
- We can use indicator random variables and linearity of expectation to compute this
Analysis

• For all $1 \leq i < j \leq m$, let $X_{i,j}$ be an indicator random variable defined such that:
  – $X_{i,j} = 1$ if person $i$ and person $j$ have the same birthday
  – $X_{i,j} = 0$ otherwise
• Note that for all $i, j$,
  
  $$E(X_{i,j}) = P(\text{person } i \text{ and } j \text{ have same birthday})$$
  
  $$= \frac{1}{n}$$
Analysis

- Let $X$ be a random variable giving the number of pairs of people with the same birthday
- We want $E(X)$
- Then $X = \sum_{1 \leq i < j \leq m} X_{i,j}$
- So $E(X) = E(\sum_{1 \leq i < j \leq m} X_{i,j})$
\[ E(X) = E \left( \sum_{1 \leq i < j \leq m} X_{i,j} \right) \]
\[ = \sum_{1 \leq i < j \leq m} E(X_{i,j}) \]
\[ = \sum_{1 \leq i < j \leq m} \frac{1}{n} \]
\[ = (m-1) \frac{1}{2n} \]
\[ = \frac{m(m-1)}{2n} \]

The second step follows by Linearity of Expectation
Reality Check

- Thus, if \( m(m - 1) \geq 2n \), expected number of pairs of people with same birthday is at least 1
- Thus if have at least \( \sqrt{2n} \) people in the room, expected number of pairs with same birthday is at least 1.
- For \( n = 365 \), if \( m = 28 \), expected number of pairs with same birthday is 1.04
In-Class Exercise

- Assume there are $m$ people in a room, and $n$ days in a year
- Assume that each of these $m$ people is born on a day chosen uniformly at random from the $n$ days
- Let $X$ be the number of groups of three people who all have the same birthday. What is $E(X)$?
- Let $X_{i,j,k}$ be an indicator r.v. which is 1 if people $i,j$, and $k$ have the same birthday and 0 otherwise
In-Class Exercise

- Q1: Write the expected value of $X$ as a function of the $X_{i,j,k}$ (use linearity of expectation)
- Q2: What is $E(X_{i,j,k})$?
- Q3: What is the total number of groups of three people out of $m$?
- Q4: What is $E(X)$?
Plan of Attack

“If you get hold of the head of a snake, the rest of it is mere rope” - Akan Proverb

- We will analyze the total number of comparisons made by quicksort
- We will let $X$ be the total number of comparisons made by R-Quicksort
- We will write $X$ as the sum of a bunch of indicator random variables
- We will use linearity of expectation to compute the expected value of $X$
Notation

- Let $A$ be the array to be sorted
- Let $z_i$ be the $i$-th smallest element in the array $A$
- Let $Z_{i,j} = \{z_i, z_i+1, \ldots, z_j\}$
Indicator Random Variables

- Let $X_{i,j}$ be 1 if $z_i$ is compared with $z_j$ and 0 otherwise.
- Note that $X = \sum_{i=1}^{n-1} \sum_{j=i+1}^{n} X_{i,j}$.
- Further note that

$$E(X) = E\left(\sum_{i=1}^{n-1} \sum_{j=i+1}^{n} X_{i,j}\right) = \sum_{i=1}^{n-1} \sum_{j=i+1}^{n} E(X_{i,j})$$
Questions

• Q1: So what is $E(X_{i,j})$?
• A1: It is $P(z_i$ is compared to $z_j$)
• Q2: What is $P(z_i$ is compared to $z_j$)?
• A2: It is:

$$P(\text{either } z_i \text{ or } z_j \text{ are the first elsms in } Z_{i,j} \text{ chosen as pivots})$$

• Why?
  - If no element in $Z_{i,j}$ has been chosen yet, no two elements in $Z_{i,j}$ have yet been compared, and all of $Z_{i,j}$ is in same list
  - If some element in $Z_{i,j}$ other than $z_i$ or $z_j$ is chosen first, $z_i$ and $z_j$ will be split into separate lists (and hence will never be compared)
More Questions

• Q: What is
  \[ P(\text{either } z_i \text{ or } z_j \text{ are first elems in } Z_{i,j} \text{ chosen as pivots}) \]

• A: \[ P(z_i \text{ chosen as first elem in } Z_{i,j}) + P(z_j \text{ chosen as first elem in } Z_{i,j}) \]

• Further note that number of elems in \( Z_{i,j} \) is \( j - i + 1 \), so
  \[ P(z_i \text{ chosen as first elem in } Z_{i,j}) = \frac{1}{j - i + 1} \]

  and
  \[ P(z_j \text{ chosen as first elem in } Z_{i,j}) = \frac{1}{j - i + 1} \]

• Hence
  \[ P(z_i \text{ or } z_j \text{ are first elems in } Z_{i,j} \text{ chosen as pivots}) = \frac{2}{j - i + 1} \]
Conclusion

\[ E(X_{i,j}) = P(z_i \text{ is compared to } z_j) \]  
\[ = \frac{2}{j - i + 1} \]
Putting it together

\[ E(X) = E \left( \sum_{i=1}^{n-1} \sum_{j=i+1}^{n} X_{i,j} \right) \] (3)

\[ = \sum_{i=1}^{n-1} \sum_{j=i+1}^{n} E(X_{i,j}) \] (4)

\[ = \sum_{i=1}^{n-1} \sum_{j=i+1}^{n} \frac{2}{j - i + 1} \] (5)

\[ = \sum_{i=1}^{n-1} \sum_{k=1}^{n-i} \frac{2}{k + 1} \] (6)

\[ < \sum_{i=1}^{n-1} \sum_{k=1}^{n} \frac{2}{k} \] (7)

\[ = \sum_{i=1}^{n-1} O(\log n) \] (8)

\[ = O(n \log n) \] (9)
Questions

• Q: Why is \( \sum_{k=1}^{n} \frac{2}{k} = O(\log n) ? \)
• A:

\[
\sum_{k=1}^{n} \frac{2}{k} = 2 \sum_{k=1}^{n} \frac{1}{k} \leq 2(\ln n + 1) \quad (10)
\]

• Where the last step follows by an integral bound on the sum (p. 1067)
How Fast Can We Sort?

• Q: What is a lowerbound on the runtime of any sorting algorithm?
• We know that $\Omega(n)$ is a trivial lowerbound
• But all the algorithms we’ve seen so far are $O(n \log n)$ (or $O(n^2)$), so is $\Omega(n \log n)$ a lowerbound?
Comparison Sorts

- Definition: An sorting algorithm is a *comparison sort* if the sorted order they determine is based only on comparisons between input elements.
- Heapsort, mergesort, quicksort, bubblesort, and insertion sort are all comparison sorts.
- We will show that any comparison sort must take $\Omega(n \log n)$. 
Comparisons

- Assume we have an input sequence $A = (a_1, a_2, \ldots, a_n)$
- In a comparison sort, we only perform tests of the form $a_i < a_j$, $a_i \leq a_j$, $a_i = a_j$, $a_i \geq a_j$, or $a_i > a_j$ to determine the relative order of all elements in $A$
- We’ll assume that all elements are distinct, and so note that the only comparison we need to make is $a_i \leq a_j$.
- This comparison gives us a yes or no answer
Decision Tree Model

- A decision tree is a full binary tree that gives the possible sequences of comparisons made for a particular input array, $A$
- Each internal node is labelled with the indices of the two elements to be compared
- Each leaf node gives a permutation of $A$
• The execution of the sorting algorithm corresponds to a path from the root node to a leaf node in the tree.
• We take the left child of the node if the comparison is $\leq$ and we take the right child if the comparison is $>$.
• The internal nodes along this path give the comparisons made by the alg, and the leaf node gives the output of the sorting algorithm.
Leaf Nodes

- Any correct sorting algorithm must be able to produce each possible permutation of the input
- Thus there must be at least $n!$ leaf nodes
- The length of the longest path from the root node to a leaf in this tree gives the worst case run time of the algorithm (i.e. the height of the tree gives the worst case runtime)
Example

- Consider the problem of sorting an array of size two: \( A = (a_1, a_2) \)
- Following is a decision tree for this problem.

```
Example

Consider the problem of sorting an array of size two: \( A = (a_1, a_2) \)
Following is a decision tree for this problem.

\[
\begin{align*}
\text{a1} \leq \text{a2?} & \\
\text{yes} & \rightarrow (a1, a2) \\
\text{no} & \rightarrow (a2, a1)
\end{align*}
\]
```
In-Class Exercise

• Give a decision tree for sorting an array of size three: \( A = (a_1, a_2, a_3) \)
• What is the height? What is the number of leaf nodes?
Height of Decision Tree

- Q: What is the height of a binary tree with at least $n!$ leaf nodes?
- A: If $h$ is the height, we know that $2^h \geq n!$
- Taking log of both sides, we get $h \geq \log(n!)$
Height of Decision Tree

• Q: What is log(n!)?
• A: It is

\[
\log(n \times (n - 1) \times \cdots \times 1) = \log n + \log(n - 1) + \cdots + \log 1 \\
\geq (n/2) \log(n/2) \\
\geq (n/2)(\log n - \log 2) \\
= \Omega(n \log n)
\]

• Thus any decision tree for sorting n elements will have a height of \(\Omega(n \log n)\)
Take Away

- We’ve just proven that any comparison-based sorting algorithm takes $\Omega(n \log n)$ time.
- This does *not* mean that *all* sorting algorithms take $\Omega(n \log n)$ time.
- In fact, there are non-comparison-based sorting algorithms which, under certain circumstances, are asymptotically faster.
Bucket Sort

- Bucket sort assumes that the input is drawn from a uniform distribution over the range \([0, 1)\)
- Basic idea is to divide the interval \([0, 1)\) into \(n\) equal size regions, or buckets
- We expect that a small number of elements in \(A\) will fall into each bucket
- To get the output, we can sort the numbers in each bucket and just output the sorted buckets in order
Bucket Sort

//PRE: A is the array to be sorted, all elements in A[i] are between 0 and 1 inclusive.
//POST: returns a list which is the elements of A in sorted order
BucketSort(A){
    B = new List[]
    n = length(A)
    for (i=1;i<=n;i++){
        insert A[i] at end of list B[floor(n*A[i])];
    }
    for (i=0;i<=n-1;i++){
        sort list B[i] with insertion sort;
    }
    return the concatenated list B[0],B[1],...,B[n-1];
}
Bucket Sort

• Claim: If the input numbers are distributed uniformly over the range $[0, 1)$, then Bucket sort takes expected time $O(n)$
• Let $T(n)$ be the run time of bucket sort on a list of size $n$
• Let $B_i$ be the random variable giving the number of elements in bucket $B[i]$
• Then $T(n) = \Theta(n) + \sum_{i=0}^{n-1} O(B_i^2)$
Analysis

- We know $T(n) = \Theta(n) + \sum_{i=0}^{n-1} O(B_i^2)$
- Taking expectation of both sides, we have

$$E(T(n)) = \Theta(n) + E\left(\sum_{i=0}^{n-1} CB_i^2\right)$$

$$= \Theta(n) + \sum_{i=0}^{n-1} E(CB_i^2)$$

$$= \Theta(n) + \sum_{i=0}^{n-1} CE(B_i^2)$$

- The second step follows by linearity of expectation
- The last step holds since for any constant $a$ and random variable $X$, $E(aX) = aE(X)$ (see Equation C.21 in the text)
• We claim that $E(B_i^2) = 2 - 1/n$
• To prove this, we define indicator random variables: $X_{ij} = 1$ if $A[j]$ falls in bucket $i$ and 0 otherwise (defined for all $i$, $0 \leq i \leq n - 1$ and $j$, $1 \leq j \leq n$)
• Thus, $B_i = \sum_{j=1}^{n} X_{ij}$
• We can now compute $E(B_i^2)$ by expanding the square and regrouping terms
Analysis

\[ E(B_i^2) = E \left( \left( \sum_{j=1}^{n} X_{ij} \right)^2 \right) \]

\[ = E \left( \sum_{j=1}^{n} \sum_{k=1}^{n} X_{ij} X_{ik} \right) \]

\[ = E \left( \sum_{j=1}^{n} X_{ij}^2 + \sum_{1 \leq j \leq n} \sum_{1 \leq k \leq n, k \neq j} X_{ij} X_{ik} \right) \]

\[ = \sum_{j=1}^{n} E \left( X_{ij}^2 \right) + \sum_{1 \leq j \leq n} \sum_{1 \leq k \leq n, k \neq j} E(X_{ij} X_{ik}) \]
Analysis

- We can evaluate the two summations separately. $X_{ij}$ is 1 with probability $1/n$ and 0 otherwise.
- Thus $E(X_{ij}^2) = 1 \times (1/n) + 0 \times (1 - 1/n) = 1/n$.
- Where $k \neq j$, the random variables $X_{ij}$ and $X_{ik}$ are independent.
- For any two independent random variables $X$ and $Y$, $E(XY) = E(X)E(Y)$ (see C.3 in the book for a proof of this).
- Thus we have that

$$E(X_{ij}X_{ik}) = E(X_{ij})E(X_{ik})$$
$$= (1/n)(1/n)$$
$$= (1/n^2)$$
Analysis

• Substituting these two expected values back into our main equation, we get:

$$E(B_i^2) = \sum_{j=1}^{n} E(X_{ij}^2) + \sum_{1\leq j \leq n} \sum_{1\leq k \leq n, k\neq j} E(X_{ij}X_{ik})$$

$$= \sum_{j=1}^{n} \frac{1}{n} + \sum_{1\leq j \leq n} \sum_{1\leq k \leq n, k\neq j} \frac{1}{n^2}$$

$$= n\left(\frac{1}{n}\right) + (n)(n - 1)(\frac{1}{n^2})$$

$$= 1 + \frac{(n - 1)}{n}$$

$$= 2 - \frac{1}{n}$$
Analysis

- Recall that $E(T(n)) = \Theta(n) + \sum_{i=0}^{n-1} O(E(B^2_i)))$
- We can now plug in the equation $E(B^2_i) = 2 - (1/n)$ to get

$$E(T(n)) = \Theta(n) + \sum_{i=0}^{n-1} 2 - (1/n)$$

$$= \Theta(n) + \Theta(n)$$

$$= \Theta(n)$$

- Thus the entire bucket sort algorithm runs in expected linear time